Optimization of Control Parameters for Adaptive Traffic-Actuated Signal Control

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ABSTRACT

This paper proposes a real-time adaptive control model for signalized intersections that decides optimal control parameters commonly found in modern actuated controllers, aiming to exploit the adaptive functionality of traffic-actuated control and to improve the performance of traffic-actuated signal system. This model incorporates a flow prediction process that estimates the future arrival rates and turning proportions at target intersection based on the available detector information and signal timing plan. Signal control parameters are optimized dynamically cycle-by-cycle to satisfy these estimated demands. The proposed adaptive control strategy is tested on a network consisting of thirty-eight actuated signals using microscopic simulation. Simulation results show the proposed adaptive model is able to improve the performance of the study network, especially under off-peak traffic conditions.

Keywords: Adaptive signal control, optimization, microscopic simulation

INTRODUCTION

At a signalized intersection, traffic signals typically operate in one of two control modes: pre-timed and actuated (semi-actuated and full-actuated). In pre-timed control, all of the control parameters, such as cycle length, phase splits and phase sequence, are preset off-line based on an assumed deterministic demand level at different time periods of day. This control mode has only limited ability to accommodate the traffic fluctuations that are commonly found in reality. In actuated control, based on the vehicle actuations registered at detectors or other traffic sensors, cycle length, phase splits and even phase sequence can vary in response to current traffic demand. However, actuated controllers change these timings subject to a set of pre-defined, fixed parameters such as unit extension and maximum green, which do not “adapt to” current conditions. Real-time adaptive signal controls offer potential to satisfy non-linear and stochastic traffic demand patterns that tax the ability of actuated control.

Existing adaptive controls that have been deployed, such as SCOOT (1) and SCATS (2, 3), which are considered “on-line” algorithms, adjust the timing parameters incrementally to accommodate changing traffic demands. This strategy is in fact to pick a “best” pre-determined off-line timing plan with detector data to “match” current traffic. The major drawback of these systems is that they are not proactive, and hence cannot satisfy significant transients effectively. RHODESTM, a real-time traffic-adaptive signal control system developed at University of Arizona, uses a traffic flow arrivals algorithm–PREDICT (4)–to improve effectiveness when calculating online phase timings. In the PREDICT algorithm, detector information on approaches of every upstream intersections, together with the traffic state, and control plan for the upstream signals are used to predict future traffic volume. It assumes that all surrounding upstream intersections have fixed-time signalized planning, an assumption that is violated in virtually every modern system.
In none of these previous systems do the embedded traffic flow prediction models fully utilize available detector information and control features. Consequently, their applicability is confined only to particular factors, and thus restricted in achieving comprehensively good performance. And, although a number of theoretical issues involving, truly adaptive, systems have been explored and proposed (5, 6, 7, 8, 9, 10, and 11), their applicability to a functional model capable of addressing real-world situations has not been tested.

For any signalized intersection, at least three kinds of information—vehicle actuated detector information, signal timing plan and current signal phase information—can be exploited to infer a relatively rich body of information that can be used in adapting the operation of a signal controller to current or expected conditions. The proposed adaptive control strategy developed in this paper incorporates a traffic flow prediction model that sufficiently utilizes all available information gleaned from signalized intersections to estimate the future arrival flows. Then optimal timing parameters are calculated based on the estimation for each signal phase in the study cycle. The optimized timing parameters are used as signal phase information for further estimations. This on-line, dynamically recursive computation process properly reflects the functionality of truly adaptive controllers. Particularly noteworthy here is that the adaptive control model proposed in this paper aims to provide an adaptive traffic-actuated control strategy for off-peak periods, and not for peak-period conditions under which fixed-time or/and actuated signal coordination control is often employed quite effectively.

This paper is organized as follows. The adaptive control model is described in detail in the next section. Then, the adaptive signal control model is tested within a microscopic simulation model, Paramics. Finally, conclusions are given.

**METHODOLOGY**

**Methodology Overview**

The proposed adaptive intersection control strategy incorporates a traffic flow prediction model that is an extension of Liu et al (12) to estimate approach volumes based on the outflows from upstream signalized intersection and forecasts turning movements at downstream intersection according to the turning fractions in previous cycles. Thus, the future flow rate associated with each phase at a downstream signal is decided by multiplying the estimated approach flow and corresponding turning proportion. Appropriate phase timing parameters are then computed in accordance with these estimated flow rates. In the model formulation, we restrict our control to only those parameters commonly found in modern actuated controllers: minimum green time, maximum green time and unit extension (gap time). Minimum green is computed based on the queue theory, determining an expected minimum green period required to dissipate the queue formed during red interval and the initial green period; maximum green is calculated according to the Webster’s function, accounting for any spillover from previous cycle. In determining gap time, we extend the approach taken by McNeil (5) to formulate a nonlinear optimization problem with objective minimizing total
intersection control delay per cycle. Using LINDO--nonlinear program solver--we obtain optimal solutions, i.e. the gap times for all phases that minimize the control delay under gap-out conditions. In any cycle in which gaps in traffic are less than this optimal gap time gap-out control does not prevail and the corresponding phase will terminate by max-out control. These optimized timing parameters are used both for the upcoming cycle as well as provide phase information in further estimations. The optimization process is shown in Figure 1.

**Figure 1 Optimization Process**

**Components of a Signal Phase**

Generally in full-actuated control, phase split $G$ consists of four portions. The first portion is the minimum green time $G_{\text{min}}$, followed by one unit extension time. If a vehicle actuates the extension detector within the unit extension, the green period is extended by another unit extension time from the moment of actuation. This fashion repeats until the termination of green interval either by max-out control, in which green extension reaches maximum allowable green, or by gap-out control, in which no vehicle is detected by extension detector during a unit extension. The second portion is defined here as the waiting time $Z$, which is the green period from the end of minimum green to the beginning of the unit extension time which incurs gap out. The third part equals unit extension $\beta$. The last part of phase split is the pre-defined, fixed change and clearance time $y+r$. The rest period in the cycle is the red duration time $R$. (Figure 2)

**Figure 2 Phase Split (phase $i$ in cycle $j$)**

Minimum green time $G_{\text{min}}$ and gap time $\beta$ are two of the objectives in our adaptive control model and will be discussed in the following subsections. However, it is worthwhile to note here that $G_{\text{min}}$ is expected to be a variant green period that is long enough to clear vehicles queued at the stop line, and $\beta$ is expected to be a particular gap time during which no vehicle actuates the extension detector, i.e. arriving vehicles during waiting time $Z$ will not be obstructed by queuing vehicles and travel freely until gap out control materializes after the specific gap time $\beta$. 

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The vehicle arrival pattern is assumed to be a Poisson process with average flow rate \( \lambda \) — strictly speaking, this assumption limits the application of the model to non-congested periods of operation. Vehicle headways then have a negative exponential distribution and the number of headways before the first one that is sufficiently large enough to invoke gap out has a geometric distribution. Then we have 

\[
\text{Probability (at least one vehicle arrives in time } \beta) = 1 - e^{-\lambda \beta}
\]

and 

\[
\text{Probability (no vehicle arrives in time } \beta) = e^{-\lambda \beta}
\]

according to the negative exponential distribution, and 

\[
\text{Probability (number of headways at which first headway larger than } \beta \text{ occurs) } = (1-e^{-\lambda \beta})^{n-1} e^{-\lambda \beta}
\]

with average value \( e^{\lambda \beta} \) according to the geometric distribution. Therefore, approximately after \( e^{\lambda \beta} \) headways, the green interval will terminate by gap-out. During these \( e^{\lambda \beta} \) headways, \( e^{\lambda \beta} \cdot 1 \) vehicles have passed the extension detector freely at flow rate \( \lambda \), then the total green extension time \( Z + \beta \) is approximately \((e^{\lambda \beta} - 1)/\lambda\). So, the waiting time \( Z \) for phase \( i \) in cycle \( j \) is

\[
Z_i^j = [\exp(\lambda_i^j \beta_i) - 1]/\lambda_i^j - \beta_i^j
\]

(1)

For ease of analysis, the displayed green interval, \( G_{\min} + Z + \beta \), is usually converted to effective green \( g_e \); then the difference of phase split and effective green is the total lost time \( L \) (start-up loss time \( l_1 \) and clearance loss time \( l_2 \), assumed constant), while the difference of cycle length and effective green is the effective red \( r_e \), which is equal to \( R + L \). (Figure 3)

When signal turns green, the first few vehicles will experience a start-up lost time \( l_1 \), after which the queued vehicles, if any, will discharge at the saturation flow rate \( S \) (assumed constant) until the queue dissipates. This green period, during which vehicles cross the stop line at saturation flow rate, is defined as queue service time or saturated green period \( P \). (Figure 4)

**Figure 3 Effective Green and Effective Red**

**Figure 4 Arrival and Discharge Pattern**
From Figure 4 and based on queue theory, the total number of vehicles departing at saturation flow during queue service time is equal to the sum of initial queue \( Q \) if any, plus arrived vehicles during effective red period and those during queue service time. Then we have

\[
S_i \times P_i^j = Q_i^j + \lambda_i^j \times (R_i^j + L_i) + \lambda_i^j \times P_i^j
\]

Therefore,

\[
P_i^j = \frac{[Q_i^j + \lambda_i^j \times (R_i^j + L_i)]}{(S_i - \lambda_i^j)}
\] (2)

It should be noted here that at the moment the queue service time expires, all queued vehicles have crossed the stop line, while it is highly possible that some following vehicles are already travelling in the road section between the stop line and the extension detector. As shown in Figure 5a for example, when queue service time expires, all queuing vehicles (in grey) have entered the intersection and the free vehicles (in white) that can enter intersection without slowing and stop, are approaching the stop line with the first one having a headway \( h \) to the stop line. Since the optimal gap time setting is expected to operate on all of the free vehicles that would traverse the intersection during green period (otherwise the calculated waiting time \( Z \) is biased), the waiting time \( Z \) is supposed to start before the end of queue service time \( P \).

Assuming that free vehicles, at least the first one, travel at a relatively constant speed from the extension detector to the stop line, we can get the situation \( U \) units of time backwards, as shown in Figure 5b. The first free vehicle is with headway \( h \) right behind the extension detector and this is the first headway on which the optimized gap time setting should operate. Note here that \( U \) is the original or the traditionally pre-set passage time based on the design speed and the set back detector’s placement. Further assuming that, when the green interval terminates, free vehicles that can use part of change and clearance interval as green extension still have \( U \) units of time to enter the intersection, we can compute the effective green and phase split as

\[
g_{ei}^j = P_i^j - U_i + Z_i^j + \beta_i^j + U_i = P_i^j + Z_i^j + \beta_i^j = \frac{[Q_i^j + \lambda_i^j \times (R_i^j + L_i)]}{(S_i - \lambda_i^j)} + \frac{\exp(\lambda_i^j \beta_i^j) - 1}{\lambda_i^j}
\] (3)

\[
G_i^j = g_{ei}^j + L_i = \frac{[Q_i^j + \lambda_i^j \times (R_i^j + S_i L_i)]}{(S_i - \lambda_i^j)} + \frac{\exp(\lambda_i^j \beta_i^j) - 1}{\lambda_i^j}
\] (4)

![Figure 5a when P expires](image1)

![Figure 5b when Z begins](image2)
Estimating Arrival Flow

In a network, the arrival flow at each leg of the downstream intersection is contributed by the corresponding phase outflows from upstream intersections, assuming no midblock exists on the links and no U-turn allowed at upstream intersections. Hence, the phase control strategy and phase timings of upstream signals are of great importance in determining the approach flow toward the downstream signal. In this paper, we introduce a flow prediction model that, first, utilizes the available phase information from upstream intersections to estimate the approach flow toward downstream intersections, and second, uses previous downstream signal phase timing to estimate the turning proportion. Referring to Figure 6 for example, two intersections are connected by a single link (right turn flow is ignored). The approach flow is the sum of outflows of phase 3 and 6 at upstream signalized intersection $k$ and consists of inflows for phase 1 and 6 at downstream signalized intersection $m$.

For upstream intersection $k$, the phase outflow is computed as the total vehicles served in the phase divided by the cycle length, then

\[ q_{ik}^j = \frac{N_{ij}}{C_{kj}} = \frac{[S_i \times P_i^j + \lambda_i^j \times (Z_i^j + \beta_i^j)]}{C_{kj}} \quad (5) \]

where

\[ C_{kj} = \text{sum} (G_{1kj} + G_{2kj} + G_{3kj} + G_{4kj}) \text{ or sum} (G_{5kj} + G_{6kj} + G_{7kj} + G_{8kj}) \quad (6) \]

For downstream intersection $m$, referring to eq.(4), we get a nonlinear function of flow rate $\lambda_i^j$, which can be solved with many methods, such as Newton’s Algorithm:

\[ F(\lambda_i^j) = \frac{[Q_i^j + \lambda_i^j R_i^j + S_i L_i]}{(S_i - \lambda_i^j ) + [\exp(\lambda_i^j \beta_i^j) - 1]/\lambda_i^j - G_i^j} = 0 \quad (7) \]

In this paper, a weighted moving average model is adopted to estimate the turning fractions $TF$, with different weights, 0.5, 0.3 and 0.2 for three previous turning proportions respectively. Then we have

\[ TF_i^j = \frac{\lambda_i^j}{(\lambda_i^j + \lambda_r^j)} \quad \text{ir} = 16, 25, 38, 47, 61, 52, 83, 74 \quad (8) \]

\[ TF_i^{j+1} = 0.5 \times TF_i^j + 0.3 \times TF_i^{j-1} + 0.2 \times TF_i^{j-2} \quad (9) \]

Therefore, the arrival rate is equal to
\[ \lambda_i^{j+1} = TF_i^{j+1} \times q_{i+r}^{j+1}, \quad \lambda_r^{j+1} = TF_r^{j+1} \times q_{i+r}^{j+1}; \quad ir = 16, 25, 38, 47, 61, 52, 83, 74 \quad (10) \]

**Determining Minimum Green**

Conventionally, minimum green is computed based on the number of vehicle actuations counted during the red period, which is a common method in volume density control type. This minimum green time, however, may not be sufficient to clear the queue formed at the stop line because there is a chance that during the minimum green time, arriving vehicles will join the queue at the stop line causing a “sluggish” congested period and not completely enter the intersection when minimum green expires. Another fact that needs to be considered is that when the minimum green time ends, the gap time setting starts operation. Considering these two factors, we set the minimum green time with introducing the queue service time mentioned previously. Then we have

\[ G_{\text{min}}^{j+1} = l_1 + P_i^{j+1} - U_i = l_1 + [Q_i^{j+1} + \lambda_i^{j+1} \times (R_i^{j+1} + L_i)] / (S_i - \lambda_i^{j+1}) - U_i \quad (11) \]

**Determining Maximum Green**

In determining maximum green, Webster’s equation

\[ D_i^{j+1} = [\lambda_i^{j+1} + Q_i^{j+1} (3600/C_{\text{max}})] / S_i \]

is used to express the effective green factor. \( C_{\text{max}} \) is the maximum allowable cycle length and is set to 120 seconds in this study. Maximum green setting is actually allocation of maximum allowable effective green time and usually consists of two computations: determine critical path and distribute green time. For both the left and right portions of the ring, we obtain the critical path and maximum effective green time as:

**Left:** Critical Path \( D_{m^*}^{j+1} + D_{n^*}^{j+1} = \max (D_1^{j+1} + D_2^{j+1}, D_5^{j+1} + D_6^{j+1}) \); \( mn=12, 56 \) (12)

**Right:** Critical Path \( D_{r^*}^{j+1} + D_{k^*}^{j+1} = \max (D_3^{j+1} + D_4^{j+1}, D_7^{j+1} + D_8^{j+1}) \); \( mn=34, 78 \) (13)

\[ G_{\text{left}}^{j+1} = [C_{\text{max}} -(L_{m^*} + L_{n^*} + L_{r^*} + L_{k^*})] (D_{m^*}^{j+1} + D_{n^*}^{j+1}) / (D_{m^*}^{j+1} + D_{n^*}^{j+1}) + (D_{r^*}^{j+1} + D_{k^*}^{j+1}) \]

(14)

\[ G_{\text{right}}^{j+1} = [C_{\text{max}} -(L_{m^*} + L_{n^*} + L_{r^*} + L_{k^*})] (D_{r^*}^{j+1} + D_{k^*}^{j+1}) / (D_{r^*}^{j+1} + D_{k^*}^{j+1}) + (D_{r^*}^{j+1} + D_{k^*}^{j+1}) \]

(15)

Then, maximum green is allocated to each phase on the critical path:

\[ G_m^{j+1} = G_{\text{left}}^{j+1} \times D_m^{j+1} / (D_{m^*}^{j+1} + D_{n^*}^{j+1}) \] (16)

\[ G_n^{j+1} = G_{\text{left}}^{j+1} \times D_n^{j+1} / (D_{m^*}^{j+1} + D_{n^*}^{j+1}) \] (17)

\[ G_r^{j+1} = G_{\text{right}}^{j+1} \times D_r^{j+1} / (D_{r^*}^{j+1} + D_{k^*}^{j+1}) \] (18)

\[ G_k^{j+1} = G_{\text{right}}^{j+1} \times D_k^{j+1} / (D_{r^*}^{j+1} + D_{k^*}^{j+1}) \] (19)

For non-critical phase, the maximum green can be calculated using the same philosophy.

Further, maximum green time \( G_{\text{max}} \) is compared with actual green time \( G-L \) or \( g_e \) to decide whether the phase terminates by gap out or max out. Then the initial queue can be determined as:

\[ \text{gap out: } Q_i^{j+1} = 0; \text{ and } \]

\[ \text{max out: } Q_i^{j+1} = \max [0, Q_i^j + \lambda_i^j \times (R_i^j + G_i) - (P_i^j S_i + \lambda_i^j Z_i^j)] \] (20)

**Determining Gap Time**

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To decide optimal gap time settings, we formulate a nonlinear optimization problem with the objective function as minimizing the total intersection control delay per cycle. The delay expression is given by Darroch (1964), which is a generalization of the well-known Webster formulation:

$$W_{ij+1} = S_i \cdot \lambda_{ij+1}^i \left[ (R_{ij+1}^i)^2 + R_{ij+1}^i (2Q_{ij+1}^i \lambda_{ij+1}^i + 1/S_i + 1/(S_i - \lambda_{ij+1}^i)) \right] / 2(S_i - \lambda_{ij+1}^i)$$  \hspace{1cm} (21)

The term $R_{ij+1}^i$ in eq.(21) can be expressed as the sum of other phases’ splits, based on the circular dependency relationship in dual-ring controller:

$$
\begin{align*}
R_{1j+1}^i &= G_{2j}^i + G_{3j}^i + G_{4j}^i \\
R_{2j+1}^i &= G_{1j+1}^i + G_{3j}^i + G_{4j}^i \\
R_{3j+1}^i &= G_{1j+1}^i + G_{2j+1}^i + G_{4j}^i \\
R_{4j+1}^i &= G_{1j+1}^i + G_{2j+1}^i + G_{3j+1}^i \\
R_{5j+1}^i &= G_{6j}^i + G_{7j}^i + G_{8j}^i \\
R_{6j+1}^i &= G_{5j+1}^i + G_{7j}^i + G_{8j}^i \\
R_{7j+1}^i &= G_{5j+1}^i + G_{6j+1}^i + G_{8j}^i \\
R_{8j+1}^i &= G_{5j+1}^i + G_{6j+1}^i + G_{7j+1}^i
\end{align*}
$$  \hspace{1cm} (22)

$G_j^i$ and $G_{ij+1}^i$ can be derived from eq.(4), therefore, in matrix form we express $R_{ij+1}^i$:

$$R_{ij+1}^i = f_i (R_j^i, \lambda_j^i, \lambda_{ij+1}^i, Q_j^i, Q_{ij+1}^i, \beta_j^i, \beta_{ij+1}^i)$$  \hspace{1cm} (23)

Substituting eq.(23) into eq.(21), we have

$$W_{ij+1}^* = \varphi_i (R_j^i, \lambda_j^i, \lambda_{ij+1}^i, Q_j^i, Q_{ij+1}^i, \beta_j^i, \beta_{ij+1}^i)$$  \hspace{1cm} (24)

Then the optimization problem is formulated as:

$$\min \sum_{i=1,2,\ldots,7,8} W_{ij+1}^i$$

subject to

$$
\begin{align*}
G_{1j+1}^i + G_{2j+1}^i &= G_{5j+1}^i + G_{6j+1}^i \\
G_{3j+1}^i + G_{4j+1}^i &= G_{7j+1}^i + G_{8j+1}^i \\
G_{1j+1}^i + G_{2j+1}^i + G_{3j+1}^i + G_{4j+1}^i &\leq C_{max} \\
G_{5j+1}^i + G_{6j+1}^i + G_{7j+1}^i + G_{8j+1}^i &\leq C_{max} \\
\beta_{ij+1}^i &\geq 0
\end{align*}
$$

The optimal gap time setting can be obtained by solving the objective function with LINDO, a nonlinear program solver. If no optimal solution is found for a certain phase, we set an arbitrary gap time for that phase, say $1/S_i$.

**TESTING**

The adaptive signal control model is tested using a scalable, high-performance microscopic simulation package, Paramics (13). Paramics has been widely used in the testing of various algorithms and evaluation of various Intelligent Transportation Systems (ITS) strategies.
because of its powerful Application Programming Interfaces (API), through which users can access the core models to customize and extend many features of the underlying simulation model without having to deal with the underlying proprietary source codes. The proposed adaptive signal control model is developed as a Paramics plugin through API programming.

The test network is shown in Figure 7, which is the so-called “Irvine Triangle” network located in southern California. A previous study calibrated the simulation network for the morning peak period from 6 to 10 AM (14). The network includes a 6-mile section of freeway I-405, a 3-mile section of freeway I-5, a 3-mile section of freeway SR-133 and several adjacent surface streets, including two parallel streets of I-405 (i.e. Alton Parkway and Barranca Parkway), one parallel streets of I-5 (Irvine Center Drive), and three crossing streets of I-405 (i.e. Culver Drive, Jeffery Road, and Sand Canyon Avenue), and etc. The network has a total of 38 actuated signals, which are operated under actuated signal coordination during peak periods and under free-mode operation during off-peak periods.

The performance of the adaptive signal control model is tested under three scenarios:

1. High demand scenario: The scenario corresponds to traffic condition in the morning peak period and its demands are obtained from the calibrated simulation model directly;
2. Medium demand scenario: its demand is equivalent to 75% of the high demand scenario;
3. Low demand scenario: its demand is equivalent to 50% of the high demand scenario.

Four performance measures are used to evaluate the performance of the proposed adaptive signal control model, which is applied to all 38 signals: (1) Average Travel Time (ATT); (2) Average Vehicle Speed (AVS); (3) Vehicle Mileage Traveled (VMT); and (4) Vehicle Hours Traveled (VHT).

Simulations are performed for each scenario for 4 hours and 15 minutes. The first 15 minutes of simulation are considered as the warm-up period for vehicles to fill in the network, and
only the last four hours of the simulations are analyzed. Five simulation runs are conducted per scenario in order to generate statistically meaningful results. The results from the median run are used for analysis and comparison. Table 1, 2 and 3 show the performances of these three scenarios.

It is found that the network under adaptive control performs better than the default actuated signal control in all three scenarios: drivers spend less time in the network, but travel more distance with improved traveling speed. It is also found that the adaptive control network performs best in scenario 2 (medium demand), probably both because the model was developed based on the assumption that vehicle arrivals are a Poisson process, and it is under precisely such conditions that gap settings control operation. On the other hand, high demand may incur large ratio of max-out controls; alternatively, low volume may result in large gap time setting. In both cases the effect of optimal gap time setting is not expected to be significant.

Table 1 Performance under low demand scenario

<table>
<thead>
<tr>
<th></th>
<th>ATT (s)</th>
<th>AVS (mph)</th>
<th>VMT (miles)</th>
<th>VHT (hours)</th>
</tr>
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<tbody>
<tr>
<td>Baseline</td>
<td>290.0</td>
<td>52.1</td>
<td>378535.85</td>
<td>7268.16</td>
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<tr>
<td>Adaptive control</td>
<td>268.6</td>
<td>56.3</td>
<td>382860.15</td>
<td>6799.97</td>
</tr>
<tr>
<td>Improvement (%)</td>
<td>-7.38</td>
<td>8.06</td>
<td>1.14</td>
<td>-6.44</td>
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Table 2 Performance under medium demand scenario

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<th>VMT (miles)</th>
<th>VHT (hours)</th>
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<tbody>
<tr>
<td>Baseline</td>
<td>340.3</td>
<td>45.2</td>
<td>557699.02</td>
<td>12534.73</td>
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<tr>
<td>Adaptive control</td>
<td>277.6</td>
<td>52.6</td>
<td>573416.69</td>
<td>10514.17</td>
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<tr>
<td>Improvement (%)</td>
<td>-18.42</td>
<td>16.38</td>
<td>2.82</td>
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Table 3 Performance under high demand scenario

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<tr>
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<td>338.4</td>
<td>44.5</td>
<td>762398.74</td>
<td>17147.97</td>
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<tr>
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<td>319.8</td>
<td>47.4</td>
<td>763687.24</td>
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<tr>
<td>Improvement (%)</td>
<td>-5.50</td>
<td>6.52</td>
<td>0.17</td>
<td>-5.91</td>
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</table>

CONCLUSIONS

This paper introduces an adaptive signal control model that can be applied to an existing actuated signal control system to improve the performance, particularly during off-peak periods. The model was tested using micro-simulation. It is found that model performance was best under traffic conditions of medium intensity, where the assumed Poisson arrival process underlying the model is expected to be valid.

In the real world, however, traffic conditions may not be consistent with the assumed Poisson process, especially at closely spaced intersections where coordinated control strategy should be taken into account. Future effort will be made to seek a more sophisticated intersection control model that includes: (1) a more accurate flow prediction model that considers platoon
and dispersion factors in estimating approach flow between linked signals; (2) optimization of control parameters for traffic flow under both high-volume and low-volume conditions.

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