Abstract. The dynamics of belief and knowledge is one of the major components of any autonomous system that should be able to incorporate new pieces of information. We introduced the Horn knowledge base dynamics to deal with two important points: first, to handle belief states that need not be deductively closed; and the second point is the ability to declare certain parts of the belief as immutable. In this paper, we address another, radically new approach to this problem. This approach is very close to the Hansson’s dyadic representation of belief. Here, we consider the immutable part as defining a new logical system. By a logical system, we mean that it defines its own consequence relation and closure operator. Based on this, we provide an abductive framework for Horn knowledge base dynamics.

Keyword: AGM, Immutable, Integrity Constraint, Knowledge Base Dynamics, Abduction.

1 Introduction

Over the last three decades [15], abduction has been embraced in AI as a non-monotonic reasoning paradigm to address some of the limitations of deductive reasoning in classical logic. The role of abduction has been demonstrated in a variety of applications. It has been proposed as a reasoning paradigm in AI for diagnosis, natural language understanding, default reasoning, planning, knowledge assimilation and belief revision, multi-agent systems and other problems (see [47]).

In the concept of knowledge assimilation and belief revision (see [38]), when a new item of information is added to a knowledge base, inconsistency can result. Revision means modifying the Horn knowledge base in order to maintain consistency, while keeping the new information and removing (contraction) or not removing the least possible previous information. In our case, update means revision and contraction, that is insertion and deletion in database perspective. Our previous work [10,11] makes connections with contraction from Horn knowledge base dynamics.

Our Horn knowledge base dynamics is defined in two parts: an immutable part (formulae or sentences) and updatable part (literals) (for definition and
properties see works of Nebel [35] and Segerberg [48]). Horn knowledge bases have a set of integrity constraints (see the definitions in later section). In the case of finite Horn knowledge bases, it is sometimes hard to see how the update relations should be modified to accomplish certain knowledge base updates.

Example 1. Consider a database with an (immutable) rule that a staff member is a person who is currently working in a research group under a chair. Additional (updatable) facts are that Matthias and Gerhard are group chairs, and Delhibabu and Aravindan are staff members in group info1. Our first integrity constraint (IC) is that each research group has only one chair. 

\[ \forall x, y, z \ (y = z) \leftarrow \text{group} \cdot \text{chair}(x, y) \land \text{group} \cdot \text{chair}(x, z). \]

Second integrity constraint is that a person can be a chair for only one research group.

\[ \forall x, y, z \ (y = z) \leftarrow \text{group} \cdot \text{chair}(y, x) \land \text{group} \cdot \text{chair}(z, x). \]

Immutable part: \( \text{staff} \cdot \text{chair}(X, Y) \leftarrow \text{staff} \cdot \text{group}(X, Z) \land \text{group} \cdot \text{chair}(Z, Y). \)

Updatable part:

\begin{align*}
\text{group} \cdot \text{chair}(\text{info1}, \text{matthias}) & \leftarrow \\
\text{group} \cdot \text{chair}(\text{info2}, \text{gerhard}) & \leftarrow \\
\text{staff} \cdot \text{group}(\text{delhibabu}, \text{info1}) & \leftarrow \\
\text{staff} \cdot \text{group}(\text{aravindan}, \text{info1}) & \leftarrow \\
\end{align*}

Suppose we want to update this database with the information, \( \text{staff} \cdot \text{chair}(\text{delhibabu}, \text{aravindan}) \), that is

\[ \text{staff} \cdot \text{chair}(\text{delhibabu}, \text{aravindan}) \leftarrow \text{staff} \cdot \text{group}(\text{delhibabu}, Z) \land \\
\text{group} \cdot \text{chair}(Z, \text{aravindan}) \]

If we are restricted to definite clauses, there is only one plausible way to do this: Delhibabu and Aravindan belong to groups info1, this updating means that we need to delete (remove) Matthias from the database and newly add (insert) Aravindan to the database (Aravindan got promoted to the chair of the research group info1 and he was removed from research group info1). This results in an update that is too strong. If we allow disjunctive information into the database, however, we can accomplish the update by minimal adding wrt consistency

\[ \text{staff} \cdot \text{group}(\text{delhibabu}, \text{info1}) \lor \text{group} \cdot \text{chair}(\text{info1}, \text{aravindan}) \]

and this option appears intuitively to be correct.

When adding new beliefs to the Horn knowledge base, if the new belief is violating integrity constraints then belief revision needs to be performed, otherwise, it is simply added. As we will see, in these cases abduction can be used in order to compute all the possibilities and it is not up to user or system to choose among them.

When dealing with the revision of a Horn knowledge base (both insertions and deletions), there are other ways to change a Horn knowledge base and it has to be performed automatically also. Considering the information, change is precious and must be preserved as much as possible. The principle of minimal
can provide a reasonable strategy. On the other hand, practical implementations have to handle contradictory, uncertain, or imprecise information, so several problems can arise: how to define efficient change in the style of AGM \([1]\), what result has to be chosen \([25,29,33]\); and finally, according to a practical point of view, what computational model to support for Horn knowledge base revision has to be provided?

Since Horn knowledge base change is one of the main problems arising in knowledge representation, it has been tackled according to several points of view. In this article, we consider the immutable part as defining a new logical system. By a logical system, we mean that it defines its own consequence relation and closure operator. Based on this, we provide an abductive framework for belief dynamics (see \([3,8,50]\)).

The rest of the paper is organized as follows: First we start with preliminaries along with the concept of logical system and properties of consequence operator. In Section 3, we introduce Horn knowledge base dynamics with our logical system. In Section 4, we explore the relationship of Horn knowledge base dynamics with coherence approach. In Section 5, we present how Horn knowledge base dynamics can be realized using abductive explanations. In Section 6, we give brief overview of related works. In Section 7, we make conclusions with a summary of our contribution as well as a discussion of future directions of investigation.

2 Preliminaries

A first-order language consists of an alphabet \(\mathcal{A}\) of a language \(\mathcal{L}\). We assume a countable universe of variables \(\text{Var}\), ranged over \(x, y, z\), and a countable universe of relation (i.e., predicate) symbols, ranged over by \(\mathcal{A}\). The following grammar defines FOL, the language of first-order logic with equality and binary relations:

\[
\phi ::= x = x | a(x, x) | \neg \phi | \bigvee \phi | \bigwedge \phi | \exists X : \phi.
\]

Here \(\phi \subseteq \text{FOL}\) and \(X \subseteq \text{Var}\) are finite sets of formulae and variables, respectively.

**Definition 1 (Normal Logic Program [22])**. By an alphabet \(\mathcal{A}\) of a language \(\mathcal{L}\) we mean disjoint sets of constants, predicate symbols, and function symbols, with at least one constant. In addition, any alphabet is assumed to contain a countably infinite set of distinguished variable symbols. A term over \(\mathcal{A}\) is defined recursively as either a variable, a constant or an expression of the form \(f(t_1, \ldots, t_n)\) where \(f\) is a function symbol of \(\mathcal{A}\), \(n\) its arity, and the \(t_i\) are terms. An atom over \(\mathcal{A}\) is an expression of the form \(P(t_1, \ldots, t_n)\) where \(P\) is a predicate symbol of \(\mathcal{A}\) and the \(t_i\) are terms. A literal is either an atom \(\mathcal{A}\) or its default negation \(\neg \mathcal{A}\). We dub default literals those of the form \(\neg \mathcal{A}\). A term (resp. atom, literal) is said ground if it does not contain variables. The set of all ground terms (resp. atoms) of \(\mathcal{A}\) is called the Herbrand universe (resp. base) of \(\mathcal{A}\). A Normal Logic Program is a possibly infinite set of rules (with no infinite descending chains of syntactical dependency) of the form:
\[ H \leftarrow B_1, ..., B_n, not C_1, ..., not C_m, \ (\text{with} \ m, n \geq 0 \ \text{and} \ \text{finit}) \]

Where \( H \), the \( B_i \) and the \( C_j \) are atoms, and each rule stands for all its ground instances. In conformity with the standard convention, we write rules of the form \( H \leftarrow \) also simply as \( H \) (known as fact). An NLP \( P \) is called definite if none of its rules contain default literals. \( H \) is the head of the rule \( r \), denoted by head\((r)\), and body\((r)\) denotes the set \( \{B_1, ..., B_n, not C_1, ..., not C_m\} \) of all the literals in the body of \( r \).

When doing problem modeling with logic programs, rules of the form

\[ \bot \leftarrow B_1, ..., B_n, not C_1, ..., not C_m, \ (\text{with} \ m, n \geq 0 \ \text{and} \ \text{finit}) \]

with a non-empty body are known as a type of integrity constraints (ICs), specifically denials, and they are normally used to prune out unwanted candidate solutions. We abuse the not default negation notation applying it to non-empty sets of literals too: we write not \( S \) to denote \( \{not s : s \in S\} \), and confound not not \( a \equiv a \). When \( S \) is an arbitrary, non-empty set of literals \( S = \{B_1, ..., B_n, not C_1, ..., not C_m\} \) we use:

- \( S^+ \) denotes the set \( \{B_1, ..., B_n\} \) of positive literals in \( S \).
- \( S^- \) denotes the set \( \{not C_1, ..., not C_m\} \) of negative literals in \( S \).
- \( |S| = S^+ \cup (not S^-) \) denotes the set \( \{B_1, ..., B_n, C_1, ..., C_m\} \) of atoms of \( S \).

As expected, we say a set of literals \( S \) is consistent iff \( S^+ \cap |S^-| = \emptyset \). We also write heads\((P)\) to denote the set of heads of non-IC rules of a (possibly constrained) program \( P \), i.e., heads\((P) = \{\text{head}(r) : r \in P\} \setminus \{\bot\} \), and facts\((P)\) to denote the set of facts of \( P \) - facts\((P) = \{\text{head}(r) : r \in P \land \text{body}(r) = \emptyset\} \).

**Definition 2 (Level mapping[4]).** Let \( P \) be a normal logic program and \( B_P \) its Herbrand base. A level mapping for \( P \) is a function \( \parallel : B_P \rightarrow \mathbb{N} \) of ground atoms to natural numbers. The mapping \( \parallel \) is also extended to ground literals by assigning \( |not A| = |A| \) for all ground atoms \( A \in B_P \). For every literal ground \( L, |L| \) is called as the level of \( L \) in \( P \).

**Definition 3 (Acyclic program [4]).** Let \( P \) be a normal logic program and \( \parallel \) a level mapping for \( P \). \( P \) is called acyclic with respect to \( \parallel \) if for every ground clause \( H \leftarrow L_1, ..., L_n \) (with \( n \geq 0 \) and \text{finit}) in \( P \) the level of \( A \) is higher then the level of every \( L_i \) (\( 1 \leq i \leq n \)). Moreover \( P \) is called acyclic if \( P \) is acyclic with respect to some level mapping for \( P \).

Unlike Horn knowledge base dynamics, where knowledge is defined as a set of sentences, here we wish to define a Horn knowledge base \( KB \) wrt a language \( \mathcal{L} \), as an abductive framework \( < P, Ab, IC, K > \), where,

* \( P \) is an acyclic normal logic program with all abducibles in \( P \) at level 0 and no non-abducible at level 0. \( P \) is referred to as a logical system. This in conjunction with the integrity constraints corresponds to immutable part of the Horn knowledge base, here \( P \) is defined by immutable part. This is discussed further in the next subsection;
* Ab is a set of atoms from $\mathcal{L}$, called the *abducibles*. This notion is required in an abductive framework, and this corresponds to the atoms that may appear in the updatable part of the knowledge;

* IC is the set of *integrity constraints*, a set of sentences from language $\mathcal{L}$. This specifies the integrity of a Horn knowledge base and forms a part of the knowledge that cannot be modified over time;

* $K$ is a set of sentences from $\mathcal{L}$. It is the *current knowledge*, and the only part of $KB$ that changes over time. This corresponds to the updatable part of the Horn knowledge base. The main requirement here is that no sentence in $K$ can have an atom that does not appear in $Ab$.

### 2.1 Logical system

The main idea of our approach is to consider the immutable part of the knowledge to define a new logical system. By a logical system, we mean that $P$ defines its own consequence relation $\models_P$ and its closure $Cn_P$. Given $P$, we have the Herbrand Base $HB_P$ and $GP$, the ground instantiation of $P$.

An *abductive interpretation* $I$ is a set of abducibles, i.e. $I \subseteq Ab$. How $I$ interprets all the ground atoms of $L$ is defined, inductively on the level of atoms wrt $P$, as follows:

* An atom $A$ at level 0 (note that only abducibles are at level 0) is interpreted as: $A$ is true in $I$ iff $A \in I$, else it is false in $I$.

* An atom $A$ at level $n$ is interpreted as: $A$ is true in $I$ iff $\exists$ clause $A \leftarrow L_1, \ldots, L_k$ in $GP$ s.t. $\forall L_j$ (1 ≤ $j$ ≤ $k$) if $L_j$ is an atom then $L_j$ is true in $I$, else if $L_j$ is a negative literal $\neg B_j$, then $B_j$ is false in $I$.

This interpretation of ground atoms can be extended, in the usual way, to interpret sentences in $L$, as follows (where $\alpha$ and $\beta$ are sentences):

* $\neg \alpha$ is true in $I$ iff $\alpha$ is false in $I$.

* $\alpha \land \beta$ is true in $I$ iff both $\alpha$ and $\beta$ are true in $I$.

* $\alpha \lor \beta$ is true in $I$ iff either $\alpha$ is true in $I$ or $\beta$ is true in $I$.

* $\forall \alpha$ is true in $I$ iff all ground instantiations of $\alpha$ are true in $I$.

* $\exists \alpha$ is true in $I$ iff some ground instantiation of $\alpha$ is true in $I$.

Given a sentence $\alpha$ in $L$, an abductive interpretation $I$ is said to be an *abductive model* of $\alpha$ iff $\alpha$ is true in $I$. Extending this to a set of sentences $K$, $I$ is an abductive model of $K$ iff $I$ is an abductive model of every sentence $\alpha$ in $K$.

Given a set of sentences $K$ and a sentence $\alpha$, $\alpha$ is said to be a $P$-consequence of $K$, written as $K \models_P \alpha$, iff every abductive model of $K$ is an abductive model of $\alpha$ also. Putting it in other words, let $Mod(K)$ be the set of all abductive models of $K$. Then $\alpha$ is a $P$-consequence of $K$ iff $\alpha$ is true in all abductive interpretations in $Mod(K)$. The consequence operator $Cn_P$ is then defined as $Cn_P(K) =$

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1 the set of all the ground atoms of $L$, in fact depends of $L$, and is given as $HB_P$, the Herbrand Base of P
\{ \alpha \mid K \models P \alpha \} = \{ \alpha \mid \alpha \text{ is true in all abductive interpretations in } \text{Mod}(K) \}.

K is said to be \textit{P-consistent} iff there is no expression \( \alpha \) s.t. \( \alpha \in \text{CnP}(K) \) and \( \neg \alpha \in \text{CnP}(K) \). Two sentences \( \alpha \) and \( \beta \) are said to be \textit{P-equivalent} to each other, written as \( \alpha \equiv \beta \), iff they have the same set of abductive models , i.e. \( \text{Mod}(\alpha) = \text{Mod}(\beta) \).

\subsection*{2.2 Properties of consequences operator}

Since a new consequence operator is defined, it is reasonable, to ask whether it satisfies certain properties that are required in the Horn knowledge base dynamics context. Here, we observe that all the required properties, listed by various researchers in Horn knowledge base dynamics, are satisfied by the defined consequence operator. The following propositions follow from the above definitions, and can be verified easily.

\( \text{CnP} \) satisfies \textit{inclusion}, i.e. \( K \subseteq \text{CnP}(K) \).

\( \text{CnP} \) satisfies \textit{iteration}, i.e. \( \text{CnP}(K) = \text{CnP}(\text{CnP}(K)) \).

Another interesting property is \textit{monotony}, i.e. if \( K \subseteq K' \), then \( \text{CnP}(K) \subseteq \text{CnP}(K') \). \( \text{CnP} \) satisfies monotony. To see this, first observe that \( \text{Mod}(K') \subseteq \text{Mod}(K) \).

\( \text{CnP} \) satisfies \textit{superclassicality} , i.e. if \( \alpha \) can be derived from \( K \) by first order classical logic, then \( \alpha \in \text{CnP}(K) \).

\( \text{CnP} \) satisfies \textit{deduction} , i.e. if \( \beta \in \text{CnP}(K \cup \{ \alpha \}) \), then \( (\beta \leftarrow \alpha) \in \text{Cn}(K) \).

\( \text{CnP} \) satisfies \textit{compactness} , i.e. if \( \alpha \in \text{CnP}(K) \), then \( \alpha \in \text{CnP}(K') \) for some finite subset \( K' \) of \( K \).

\subsection*{2.3 Statics of a Horn knowledge base}

The statics of a Horn knowledge base \( KB \), is given by the current knowledge \( K \) and the integrity constraints \( IC \). An abductive interpretation \( M \) is an abductive model of \( KB \) iff it is an abductive model of \( K \cup IC \). Let \( \text{Mod}(KB) \) be the set of all abductive models of \( KB \). The belief set represented by \( KB \), written as \( KB^* \) is given as,

\[ KB^* = \text{CnP}(K \cup IC) = \{ \alpha \mid \alpha \text{ is true in every abductive model of } KB \}. \]

A belief (represented by a sentence in \( L \)) \( \alpha \) is \textit{accepted} in \( KB \) iff \( \alpha \in KB^* \) (i.e. \( \alpha \) is true in every model of \( KB \)). \( \alpha \) is \textit{rejected} in \( KB \) iff \( \neg \alpha \in KB^* \) (i.e. \( \alpha \) is false in every model of \( KB \)). Note that there may exist a sentence \( \alpha \) s.t. \( \alpha \) is neither accepted nor rejected in \( KB \) (i.e. \( \alpha \) is true in some but not all models of \( KB \)), and so \( KB \) represents a partial description of the world.

Two Horn knowledge bases \( KB_1 \) and \( KB_2 \) are said to be \textit{equivalent} to each other, written as \( KB_1 \equiv KB_2 \), iff they are based on the same logical system and their current knowledge are \( P \)-equivalent, i.e. \( P_1 = P_2 \), \( Ab_1 = Ab_2 \), \( IC_1 = IC_2 \) and \( K_1 \equiv K_2 \). Obviously, two equivalent Horn knowledge bases \( KB_1 \) and \( KB_2 \) represent the same belief set, i.e. \( KB_1^* = KB_2^* \).
3 Horn knowledge base dynamics

In AGM [1] three kinds of belief dynamics are defined: expansion, contraction and revision. We consider all of them, one by one, in the sequel.

3.1 Expansion

Let \( \alpha \) be new information that has to be added to a knowledge base \( KB \). Suppose \( \neg \alpha \) is not accepted in \( KB \). Then, obviously \( \alpha \) is \( P \) - consistent with \( IC \), and \( KB \) can be expanded by \( \alpha \), by modifying \( K \) as follows:

\[
KB + \alpha \equiv < P, Ab, IC, K \cup \{ \alpha \} >
\]

Note that we do not force the presence of \( \alpha \) in the new \( K \), but only say that \( \alpha \) must be in the belief set represented by the expanded Horn knowledge base. If in case \( \neg \alpha \) is accepted in \( KB \) (in other words, \( \alpha \) is inconsistent with \( IC \)), then expansion of \( KB \) by \( \alpha \) results in a inconsistent Horn knowledge base with no abductive models, i.e. \((KB + \alpha)^*\) is the set of all sentences in \( L \).

Putting it in model-theoretic terms, \( KB \) can be expanded by a sentence \( \alpha \), when \( \alpha \) is not false in all models of \( KB \). The expansion is defined as:

\[
Mod(KB + \alpha) = Mod(KB) \cap Mod(\alpha).
\]

If \( \alpha \) is false in all models of \( KB \), then clearly \( Mod(KB+\alpha) \) is empty, implying that expanded Horn knowledge base is inconsistent.

3.2 Revision

As usual, for revising and contracting a Horn knowledge base, the rationality of the change is discussed first. Later a construction is provided that complies with the proposed rationality postulates.
Rationality postulates

Let $KB = < P, Ab, IC, K >$ be revised by a sentence $\alpha$ to result in a new Horn knowledge base $KB + \alpha = < P', Ab', IC', K' >$.

When a Horn knowledge base is revised, we do not (generally) wish to modify the underlying logical system $P$ or the set of abducibles $Ab$. This is referred to as \textit{inferential constancy} by Hansson [19,20].

\begin{itemize}
  \item[(+1)] \textit{(Inferential constancy)} $P' = P$ and $Ab' = Ab, IC' = IC$.
  \item[(+2)] \textit{(Success)} $\alpha$ is accepted in $KB + \alpha$, i.e. $\alpha$ is true in all models of $KB + \alpha$.
  \item[(+3)] \textit{(Consistency)} $\alpha$ is satisfiable and $P$-consistent with IC iff $KB + \alpha$ is $P$-consistent, i.e. $\text{Mod}(\{\alpha\} \cup IC)$ is not empty iff $\text{Mod}(KB + \alpha)$ is not empty.
  \item[(+4)] \textit{(Vacuity)} If $\neg \alpha$ is not accepted in KB, then $KB + \alpha \equiv KB + \alpha$, i.e. if $\alpha$ is not false in all models of KB, then $\text{Mod}(KB + \alpha) = \text{Mod}(KB) \cap \text{Mod}(\alpha)$.
  \item[(+5)] \textit{(Preservation)} If $KB \equiv KB'$ and $\alpha \equiv \beta$, then $KB + \alpha \equiv KB' + \beta$, i.e. if $\text{Mod}(KB) = \text{Mod}(KB')$ and $\text{Mod}(\alpha) = \text{Mod}(\beta)$, then $\text{Mod}(KB + \alpha) = \text{Mod}(KB + \beta)$.
  \item[(+6)] \textit{(Extended Vacuity 1)} $(KB + \alpha) + \beta$ implies $KB + (\alpha \land \beta)$, i.e. $(\text{Mod}(KB + \alpha) \cap \text{Mod}(\beta)) \subseteq \text{Mod}(KB + (\alpha \land \beta))$.
  \item[(+7)] \textit{(Extended Vacuity 2)} If $\neg \beta$ is not accepted in $(KB + \alpha)$, then $(KB + (\alpha \land \beta))$ implies $(KB + \alpha) + \beta$, i.e. if $\beta$ is not false in all models of $KB + \alpha$, then $\text{Mod}(KB + (\alpha \land \beta)) \subseteq (\text{Mod}(KB + \alpha) \cap \text{Mod}(\beta))$.
\end{itemize}

Construction

Let $S$ stand for the set of all abductive interpretations that are consistent with $IC$, i.e. $S = \text{Mod}(IC)$. We do not consider abductive interpretations that are not models of $IC$, simply because $IC$ does not change during revision. Observe that when $IC$ is empty, $S$ is the set of all abductive interpretations. Given a Horn knowledge base $KB$, and two abductive interpretations $I_1$ and $I_2$ from $S$, we can compare how close these interpretations are to $KB$ by using an order $\leq_{KB}$ among abductive interpretations in $S$. $I_1 <_{KB} I_2$ iff $I_1 \leq_{KB} I_2$ and $I_2 \not\leq_{KB} I_1$.

Let $\mathcal{F} \subseteq S$. An abductive interpretation $I \in \mathcal{F}$ is minimal in $\mathcal{F}$ wrt $\leq_{KB}$ if there is no $I' \in \mathcal{F}$ s.t. $I' <_{KB} I$. Let, $\text{Min}(\mathcal{F}, \leq_{KB}) = \{I \mid I \text{ is minimal in } \mathcal{F} \text{ wrt } \leq_{KB}\}$.

For any Horn knowledge base $KB$, the following are desired properties of $\leq_{KB}$:

\begin{itemize}
  \item[$(\leq 1)$] \textit{(Pre-order)} $\leq_{KB}$ is a pre-order, i.e. it is transitive and reflexive.
  \item[$(\leq 2)$] \textit{(Connectivity)} $\leq_{KB}$ is total in $S$, i.e. \forall $I_1, I_2 \in S$: either $I_1 \leq_{KB} I_2$ or $I_2 \leq_{KB} I_1$.
  \item[$(\leq 3)$] \textit{(Faithfulness)} $\leq_{KB}$ is faithful to $KB$, i.e. $I \in \text{Min}(S, \leq_{KB})$ iff $I \in \text{Mod}(KB)$.
  \item[$(\leq 4)$] \textit{(Minimality)} For any non-empty subset $\mathcal{F}$ of $S$, $\text{Min}(\mathcal{F}, \leq_{KB})$ is not empty.
  \item[$(\leq 5)$] \textit{(Preservance)} For any Horn knowledge base $KB'$, if $KB \equiv KB'$ then $\leq_{KB} = \leq_{KB'}$.
\end{itemize}
Let $KB$ (and consequently $K$) be revised by a sentence $\alpha$, and $\leq_{KB}$ be a rational order that satisfies ($\leq 1$) to ($\leq 5$). Then the abductive models of the revised Horn knowledge base are given precisely by: $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$. Note that, this construction does not say what the resulting $K$ is, but merely says what should be the abductive models of the new Horn knowledge base.

**Representation theorem**

Now, we proceed to show that revision of $KB$ by $\alpha$, as constructed above, satisfies all the rationality postulates stipulated in the beginning of this section. This is formalized by the following lemma.

**Lemma 1.** Let $KB$ be a Horn knowledge base, $\leq_{KB}$ an order among $S$ that satisfies ($\leq 1$) to ($\leq 5$). Let a revision operator $\downarrow$ be defined as: for any sentence $\alpha$, $\text{Mod}(KB \downarrow \alpha) = \text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$. Then $\downarrow$ satisfies all the rationality postulates for revision ($\downarrow 1$) to ($\downarrow 7$).

**Proof.**

($\downarrow 1$) $P' = P$ and $Ab' = Ab$ and $IC' = IC$

This is satisfied obviously, since our construction does not touch $P$ and $Ab$, and $IC$ follows from every abductive interpretation in $\text{Mod}(KB + \alpha)$.

($\downarrow 2$) $\alpha$ is accepted in $KB + \alpha$

Note that every abductive interpretation $M \in \text{Mod}(KB + \alpha)$ is a model of $\alpha$. Hence $\alpha$ is accepted in $KB + \alpha$.

($\downarrow 3$) $\alpha$ is satisfiable and $P$-consistent with $IC$ iff $KB + \alpha$ is $P$-consistent.

If part: If $KB + \alpha$ is $P$-consistent, then $\text{Mod}(KB + \alpha)$ is not empty. This implies that $\text{Mod}(\{\alpha\} \cup IC)$ is not empty, and hence $\alpha$ is satisfiable and $P$-consistent with $IC$.

Only if part: If $\alpha$ is satisfiable and $P$-consistent with $IC$, then $\text{Mod}(\{\alpha\} \cup IC)$ is not empty, and ($\leq 4$) ensures that $\text{Mod}(KB + \alpha)$ is not empty. Thus, $KB + \alpha$ is $P$-consistent.

($\downarrow 4$) If $\neg \alpha$ is not accepted in $KB$, then $KB + \alpha \equiv KB + \alpha$.

We have to establish that $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB}) = \text{Mod}(KB) \cap \text{Mod}(\alpha)$.

Since $\neg \alpha$ is not accepted in $KB$, $\text{Mod}(KB) \cap \text{Mod}(\alpha)$ is not empty. The required result follows immediately from the fact that $\leq_{KB}$ is faithful to $KB$ (i.e. satisfies $\leq 3$), which selects only and all those models of $\alpha$ which are also models of $KB$.

($\downarrow 5$) If $KB \equiv KB'$ and $\alpha \equiv \beta$ then $KB + \alpha = KB' + \beta$ ($\leq 5$) ensures that $\leq_{KB} = \leq_{KB'}$. The required result follows immediately from this and the fact that $\text{Mod}(\alpha) = \text{Mod}(\beta)$.

($\downarrow 6$) $(KB + \alpha) + \beta$ implies $KB + (\alpha \land \beta)$.

We consider this in two cases. When $\neg \beta$ is accepted in $KB + \alpha$, $(KB + \alpha) + \beta$ is the set of all sentences from $L$, and the postulate follows immediately. Instead when $\neg \beta$ is not accepted in $KB + \alpha$, this postulates coincides with the next one.
If $\neg \beta$ is not accepted in $KB + \alpha$, then $KB + (\alpha \land \beta)$ implies $(KB + \alpha) + \beta$. Together with the second case of previous postulate, we need to show that $KB + (\alpha \land \beta) = (KB + \alpha) + \beta$. In other words, we have to establish that $\text{Min}(\text{Mod}(\{\alpha \land \beta\} \cup IC), \leq_{KB}) = \text{Mod}(KB + \alpha) \cap \text{Mod}(\beta)$. For the sake of simplicity, let us represent $\text{Min}(\text{Mod}(\{\alpha \land \beta\} \cup IC), \leq_{KB})$ by $P$, and $\text{Mod}(KB + \alpha) \cap \text{Mod}(\beta)$, which is the same as $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB}) \cap \text{Mod}(\beta)$, by $Q$. The required result is obtained in two parts:

1) $\forall$ (abductive interpretation) $M$: if $M \in P$, then $M \in Q$
   
   Obviously $M \in \text{Mod}(\beta)$. Assume that $M \notin \text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$.
   
   This can happen in two cases, and we show that both the cases lead to contradiction.
   
   Case A: No model of $\beta$ is selected by $\leq_{KB}$ from $\text{Mod}(\{\alpha\} \cup IC)$. But this contradicts our initial condition that $\neg \beta$ is not accepted in $KB + \alpha$.
   
   Case B: Some model, say $M'$, of $\beta$ is selected by $\leq_{KB}$ from $\text{Mod}(\{\alpha\} \cup IC)$. Since $M$ is not selected, it follows that $M' <_{KB} M$. But then this contradicts our initial assumption that $M \in P$. So, $P \subseteq Q$.

2) $\forall$ (abductive interpretation) $M$: if $M \in Q$, then $M \in P$
   
   $M \in Q$ implies that $M$ is a model of both $\alpha$ and $\beta$, and $M$ is selected by $\leq_{KB}$ from $\text{Mod}(\{\alpha\} \cup IC)$. Note that $\text{Mod}(\{\alpha \land \beta\} \cup IC) \subseteq \text{Mod}(\{\alpha\} \cup IC)$. Since $M$ is selected by $\leq_{KB}$ in a bigger set (i.e. $\text{Mod}(\{\alpha\} \cup IC)$), $\leq_{KB}$ must select $M$ from its subset $\text{Mod}(\{\alpha \land \beta\} \cup IC)$ also. Hence $Q \subseteq P$.

But, that is not all. Any rational revision of $KB$ by $\alpha$, that satisfies all the rationality postulates, can be constructed by our construction method, and this is formalized below.

**Lemma 2.** Let $KB$ be a Horn knowledge base and $+$ a revision operator that satisfies all the rationality postulates for revision $(\dagger 1)$ to $(\dagger 7)$. Then, there exists an order $\leq_{KB}$ among $S$, that satisfies $(\leq 1)$ to $(\leq 5)$, and for any sentence $\alpha$, $\text{Mod}(KB + \alpha)$ is given in $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$.

**Proof.** Let us construct an order $\leq_{KB}$ among interpretations in $S$ as follows:

For any two abductive interpretations $I$ and $I'$ in $S$, define $I \leq_{KB} I'$ iff either $I \in \text{Mod}(KB)$ or $I \in \text{Mod}(KB + \text{form}(I, I'))$, where $\text{form}(I, I')$ stands for sentence whose only models are $I$ and $I'$. We will show that $\leq_{KB}$ thus constructed satisfies $(\leq 1)$ to $(\leq 5)$ and $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB}) = \text{Mod}(KB + \alpha)$.

First, we show that $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB}) = \text{Mod}(KB + \alpha)$. Suppose $\alpha$ is not satisfiable, i.e. $\text{Mod}(\alpha)$ is empty, or $\alpha$ does not satisfy $IC$, then there are no abductive models of $\{\alpha\} \cup IC$, and hence $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$ is empty. From $(\dagger 3)$, we infer that $\text{Mod}(KB + \alpha)$ is also empty. When $\alpha$ is satisfiable and $\alpha$ satisfies $IC$, the required result is obtained in two parts:

1) If $I \in \text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$, then $I \in \text{Mod}(KB + \alpha)$
   
   Since $\alpha$ is satisfiable and consistent with $IC$, $(\dagger 3)$ implies that there exists at least one model, say $I'$, for $KB + \alpha$. From $(\dagger 1)$, it is clear that $I'$ is a model of $IC$, from $(\dagger 2)$ we also get that $I'$ is a model of $\alpha$, and consequently $I \leq_{KB} I'$
(because $I \in \text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$). Suppose $I \in \text{Mod}(KB)$, then (4) immediately gives $I \in \text{Mod}(KB + \alpha)$. If not, from our definition of $\leq_{KB}$, it is clear that $I \in \text{Mod}(KB + \text{form}(I, I'))$. Note that $\alpha \wedge \text{form}(I, I') = \text{form}(I, I')$, since both $I$ and $I'$ are models of $\alpha$. From (6) and (7), we get $\text{Mod}(KB + \alpha) \cap \{I, I'\} = \text{Mod}(KB + \text{form}(I, I'))$. Since $I \in \text{Mod}(KB + \text{form}(I, I'))$, it immediately follows that $I \in \text{Mod}(KB + \alpha)$.

2) If $I \in \text{Mod}(KB + \alpha)$, then $I \in \text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$.

From (1) we get $I$ is a model of $IC$, and from (2), we obtain $I \in \text{Mod}(\alpha)$. Suppose $I \in \text{Mod}(KB)$, then from our definition of $\leq_{KB}$, we get $I \leq_{KB} I'$, for any other model $I'$ of $\alpha$ and $IC$, and hence $I \in \text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$. Instead, if $I$ is not a model of $KB$, then, to get the required result, we should show that $I \in \text{Mod}(KB + \text{form}(I, I'))$, for every model $I'$ of $\alpha$ and $IC$. As we have observed previously, from (6) and (7), we get $\text{Mod}(KB + \alpha) \cap \{I, I'\} = \text{Mod}(KB + \text{form}(I, I'))$. Since $I \in \text{Mod}(KB + \alpha)$, it immediately follows that $I \in \text{Mod}(KB + \text{form}(I, I'))$. Hence $I \leq_{KB} I'$ for any model $I'$ of $\alpha$ and $IC$, and consequently, $I \in \text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$.

Now we proceed to show that the order $\leq_{KB}$ among $S$, constructed as per our definition, satisfies all the rationality axioms (1) to (5).

\[(\leq 1)\] $\leq_{KB}$ is a pre-order.

Note that we need to consider only abductive interpretations from $S$. From (2) and (3), we have $\text{Mod}(KB + \text{form}(I, I')) = \{I\}$, and so $I \leq_{KB} I$. Thus $\leq_{KB}$ satisfies reflexivity.

Let $I_1 \in \text{Mod}(IC)$ and $I_2 \notin \text{Mod}(IC)$. Clearly, it is possible that two interpretations $I_1$ and $I_2$ are not models of $KB$, and $\text{Mod}(KB + \text{form}(I_1, I_2)) = \{I_1\}$. So, $I_1 \leq_{KB} I_2$ does not necessarily imply $I_2 \leq_{KB} I_1$, and thus $\leq_{KB}$ satisfies anti-symmetry.

To show the transitivity, we have to prove that $I_1 \leq_{KB} I_3$, when $I_1 \leq_{KB} I_2$ and $I_2 \leq_{KB} I_3$ hold. Suppose $I_1 \in \text{Mod}(KB)$, then $I_1 \leq_{KB} I_3$ follows immediately from our definition of $\leq_{KB}$. On the other case, when $I_1 \notin \text{Mod}(KB)$, we first observe that $I_1 \in \text{Mod}(KB + \text{form}(I_1, I_2))$, which follows from definition of $\leq_{KB}$ and $I_1 \leq_{KB} I_2$. Also observe that $I_2 \notin \text{Mod}(KB)$. If $I_2$ were a model of $KB$, then it follows from (4) that $\text{Mod}(KB + \text{form}(I_1, I_2)) = \text{Mod}(KB) \cap \{I_1, I_2\} = \{I_2\}$, which is a contradiction, and so $I_2 \notin \text{Mod}(KB)$. This, together with $I_2 \leq_{KB} I_3$, implies that $I_2 \in \text{Mod}(KB + \text{form}(I_2, I_3))$.

Now consider $\text{Mod}(KB + \text{form}(I_1, I_2, I_3))$. Since $+$ satisfies (2) and (3), it follows that this is a non-empty subset of $\{I_1, I_2, I_3\}$. We claim that $\text{Mod}(KB + \text{form}(I_1, I_2, I_3)) \cap \{I_1, I_2\}$ can not be empty. If it is empty, then it means that $\text{Mod}(KB + \text{form}(I_1, I_2, I_3)) = \{I_3\}$. Since $+$ satisfies (6) and (7), this further implies that $\text{Mod}(KB + \text{form}(I_2, I_3)) = \text{Mod}(KB + \text{form}(I_1, I_2, I_3)) \cap \{I_2, I_3\} = \{I_3\}$. This contradicts our observation that $I_2 \in \text{Mod}(KB + \text{form}(I_2, I_3))$, and so $\text{Mod}(KB + \text{form}(I_1, I_2, I_3)) \cap \{I_1, I_2\}$ can not be empty. Using (6) and (7) again, we get $\text{Mod}(KB + \text{form}(I_1, I_2)) = \text{Mod}(KB + \text{form}(I_1, I_2, I_3)) \cap \{I_1, I_2\}$. Since we know that $I_1 \in \text{Mod}(KB + \text{form}(I_1, I_2))$, it follows that $I_1 \in \text{Mod}(KB + \text{form}(I_1, I_2, I_3))$. From (6)
and (\(\dagger\dagger\)) we also get \(\text{Mod}(KB + \text{form}(I_1, I_3)) = \text{Mod}(KB + \text{form}(I_1, I_2, I_3))\cap \{I_1, I_3\}\), which clearly implies that \(I_1 \in \text{Mod}(KB + \text{form}(I_1, I_3))\). From our definition of \(\leq_{KB}\), we now obtain \(I_1 \leq_{KB} I_3\). Thus, \(\leq_{KB}\) is a pre-order.

\((\leq 2)\) \(\leq_{KB}\) is total.

Since + satisfies (\(\dagger\dagger\)) and (\(\dagger\dagger\)), for any two abductive interpretations \(I\) and \(I'\) in \(S\), it follows that \(\text{Mod}(KB + \text{form}(I, I'))\) is a non-empty subset of \(\{I, I'\}\). Hence, \(\leq_{KB}\) is total.

\((\leq 3)\) \(\leq_{KB}\) is faithful to \(KB\).

From our definition of \(\leq_{KB}\), it follows that \(\forall I_1, I_2 \in \text{Mod}(KB) : I_1 <_{KB} I_2\) does not hold. Suppose \(I_1 \in \text{Mod}(KB)\) and \(I_2 \notin \text{Mod}(KB)\). Then, we have \(I_1 \leq_{KB} I_2\). Since + satisfies (\(\dagger\dagger\)), we also have \(\text{Mod}(KB + \text{form}(I_1, I_2)) = \{I_1\}\). Thus, from our definition of \(\leq_{KB}\), we can not have \(I_2 \leq_{KB} I_1\). So, if \(I_1 \in \text{Mod}(KB)\) and \(I_2 \notin \text{Mod}(KB)\), then \(I_1 <_{KB} I_2\) holds. Thus, \(\leq_{KB}\) is faithful to \(KB\).

\((\leq 4)\) For any non-empty subset \(F\) of \(S\), \(\text{Min}(F, \leq_{KB})\) is not empty.

Let \(\alpha\) be a sentence such that \(\text{Mod}((\{\alpha\} \cup IC) = F\). We have already shown that \(\text{Mod}(KB + \alpha) = \text{Min}(F, \leq_{KB})\). Since, + satisfies (\(\dagger\dagger\)), it follows that \(\text{Mod}(KB + \alpha)\) is not empty, and thus \(\text{Min}(F, \leq_{KB})\) is not empty.

\((\leq 5)\) If \(KB \equiv KB'\), then \(\leq_{KB} = \leq_{KB'}\).

This follows immediately from the fact that + satisfies (\(\dagger\dagger\)).

Thus, the order among interpretations \(\leq_{KB}\), constructed as per our definition, satisfies (\(\leq 1\)) to (\(\leq 5\)), and \(\text{Mod}(KB + \alpha) = \text{Min}(\text{Mod}((\{\alpha\} \cup IC), \leq_{KB})\).

So, we have a one to one correspondence between the axiomatization and the construction, which is highly desirable, and this is summarized by the following representation theorem.

**Theorem 1.** Let \(KB\) be revised by \(\alpha\), and \(KB + \alpha\) be obtained by the construction discussed above. Then, + is a revision operator iff it satisfies all the rationality postulates (\(+1\)) to (\(+7\)).

**Proof.** Follows from Lemma 1. and Lemma 2. ■

### 3.3 Contraction

Contraction of a sentence from a Horn knowledge base \(KB\) is studied in the same way as that of revision. We first discuss the rationality of change during contraction and proceed to provide a construction for contraction using duality between revision and contraction.

**Rationality Postulates**

Let \(KB = \langle P, Ab, IC, K \rangle\) be contracted by a sentence \(\alpha\) to result in a new Horn knowledge base \(KB - \alpha = \langle P', Ab', IC', K' \rangle\).

\((-1)\) (Inferential Constancy) \(P' = P\) and \(Ab' = Ab\) and \(IC' = IC\).
(−2) *(Success)* If \( \alpha \notin Cn_P(KB) \), then \( \alpha \) is not accepted in \( KB^\prime - \alpha \), i.e. if \( \alpha \) is not true in all the abductive interpretations, then \( \alpha \) is not true in all abductive interpretations in \( Mod(KB^\prime - \alpha) \).

(−3) *(Inclusion)* \( \forall (\text{belief}) \beta: \beta \text{ is accepted in } KB^\prime - \alpha, \text{ then } \beta \text{ is accepted in } KB \), i.e. \( Mod(KB) \subseteq Mod(KB^\prime - \alpha) \).

(−4) *(Vacuity)* If \( \alpha \) is not accepted in \( KB \), then \( KB^\prime - \alpha = KB \), i.e. if \( \alpha \) is not true in all the abductive models of \( KB \), then \( Mod(KB^\prime - \alpha) = Mod(KB) \).

(−5) *(Recovery)* \((KB^\prime - \alpha)^{+}\) implies \( KB \), i.e. \( Mod(KB^\prime - \alpha) \cap Mod(\alpha) \subseteq Mod(KB) \).

(−6) *(Preservation)* If \( KB \equiv KB^\prime \) and \( \alpha \equiv \beta \), then \( KB^\prime - \alpha = KB^\prime - \beta \), i.e. if \( Mod(KB) = Mod(KB^\prime) \) and \( Mod(\alpha) = Mod(\beta) \), then \( Mod(KB^\prime - \alpha) = Mod(KB^\prime - \beta) \).

(−7) *(Conjunction 1)* \( KB^\prime - (\alpha \wedge \beta) \) implies \( KB^\prime - \alpha \cap KB^\prime - \beta \), i.e. \( Mod(KB^\prime - (\alpha \wedge \beta)) \subseteq Mod(KB^\prime - \alpha) \cup Mod(KB^\prime - \beta) \).

(−8) *(Conjunction 2)* If \( \alpha \) is not accepted in \( KB^\prime - (\alpha \wedge \beta) \), then \( KB^\prime - \alpha \) implies \( KB^\prime - (\alpha \wedge \beta) \), i.e. if \( \alpha \) is not true in all the models of \( KB^\prime - (\alpha \wedge \beta) \), then \( Mod(KB^\prime - \alpha) \subseteq Mod(KB^\prime - (\alpha \wedge \beta)) \).

Before providing a construction for contraction, we wish to study the duality between revision and contraction. The Levi and Harper identities still hold in our case, and is discussed in the sequel.

**Relationship between contraction and revision**

Contraction and revision are related to each other. Given a contraction function \( \frown \), a revision function \( \frown \) can be obtained as follows:

\[
\text{(Levi Identity)} \quad Mod(KB \frown \alpha) = Mod(KB^\prime - \alpha) \cap Mod(\alpha)
\]

The following theorem formally states that Levi identity holds in our approach.

**Theorem 2.** Let \( \frown \) be a contraction operator that satisfies all the rationality postulates (−1) to (−8). Then, the revision function \( \frown \), obtained from \( \frown \) using the Levi Identity, satisfies all the rationality postulates (−1) to (−7).

Similarly, a contraction function \( \frown \) can be constructed using the given revision function \( \frown \) as follows:

\[
\text{(Harper Identity)} \quad Mod(KB^\prime - \alpha) = Mod(KB) \cup Mod(KB \frown - \alpha)
\]

**Theorem 3.** Let \( \frown \) be a revision operator that satisfies all the rationality postulates (−1) to (−7). Then, the contraction function \( \frown \), obtained from \( \frown \) using the Harper Identity, satisfies all the rationality postulates (−1) to (−8).

**Construction**

Given the construction for revision, based on order among interpretation in \( S \), a construction for contraction can be provided as:

\[
Mod(KB^\prime - \alpha) = Mod(KB) \cup Min(Mod(\{\sim \alpha \} \cup IC), \leq KB),
\]
where $\leq_{KB}$ is the relation among interpretations in $\mathcal{S}$ that satisfies the rationality axioms ($\leq 1$) to ($\leq 5$). As in the case of revision, this construction says what should be the models of the resulting Horn knowledge base, and does not explicitly say what the resulting Horn knowledge base is.
Representation theorem

Since the construction for contraction is based on a rational contraction for revision, the following lemmae and theorem follow obviously.

**Lemma 3.** Let $KB$ be a Horn knowledge base, $\leq_{KB}$ an order among $S$ that satisfies $(\leq 1)$ to $(\leq 5)$. Let a contraction operator $\hat{-}$ be defined as: for any sentence $\alpha$, $\text{Mod}(KB\hat{-}\alpha) = \text{Mod}(KB) \cup \text{Min}(\text{Mod}(\{\neg\alpha\} \cup IC), \leq_{KB})$. Then $\hat{-}$ satisfies all the rationality postulates for contraction $(\hat{-}1)$ to $(\hat{-}8)$.

**Proof.** Follows from Theorem 1 and Theorem 3. ■.

**Lemma 4.** Let $KB$ be a Horn knowledge base and $\hat{-}$ a contraction operator that satisfies all the rationality postulates for contraction $(\hat{-}1)$ to $(\hat{-}8)$. Then, there exists an order $\leq_{KB}$ among $S$, that satisfies$(\leq 1)$ to $(\leq 5)$, and for any sentence $\alpha$, $\text{Mod}(KB\hat{-}\alpha)$ is given as $\text{Mod}(KB) \cup \text{Min}(\text{Mod}(\{\neg\alpha\} \cup IC), \leq_{KB})$.

**Proof.** Follows from Theorem 1 and Theorem 3. ■

**Theorem 4.** Let $KB$ be contracted by $\alpha$, and $KB\hat{-}\alpha$ be obtained by the construction discussed above. Then $\hat{-}$ is a contraction operator iff it satisfies all the rationality postulates $(\hat{-}1)$ to $(\hat{-}8)$.

**Proof.** Follows from Lemma 3 and Lemma 4. ■

4 Relationship with the coherence approach of AGM

Given Horn knowledge base $KB = (P, Ab, IC, K)$ represents a belief set $KB^*$ that is closed under $CnP$. We have defined how $KB$ can be expanded, revised, or contracted. The question now is: does our foundational approach (wrt classical first-order logic) on $KB$ coincide with coherence approach (wrt our consequence operator $CnP$) of AGM on $KB^*$? There is a problem in answering this question (similar practical problem [5]) , since our approach, we require $IC$ to be immutable, and only the current knowledge $K$ is allowed to change. On the contrary, AGM approach treat every sentence in $KB^*$ equally, and can throw out sentences from $CnP(IC)$. One way to solve this problem is to assume that sentences in $CnP(IC)$ are more entrenched than others. However, one-to-one correspondence can be established, when $IC$ is empty. The key is our consequence operator $CnP$, and in the following, we show that coherence approach of AGM with this consequence operator, is exactly same as our foundational approach, when $IC$ is empty.

4.1 Expansion

Expansion in AGM (see [Π]- framework is defined as $KB\#\alpha = CnP(KB^* \cup \{\alpha\})$, is is easy to see that this is equivalent to our definition of expansion (when $IC$ is empty), and is formalized below.
Theorem 5. Let $KB + \alpha$ be an expansion of $KB$ by $\alpha$ (as defined in section 3.2). Then $(KB + \alpha)^* = KB\#\alpha$.

Proof. By our definition of expansion, $(KB + \alpha)^* = Cn_P(\mathcal{I}C \cup K \cup \{\alpha\})$, which is clearly the same set as $Cn_P(KB^* \cup \{\alpha\})$. ■

4.2 Revision

AGM puts forward rationality postulates (*1) to (*8) to be satisfied by a revision operator on $KB^*$. reproduced below:

(*1) (Closure) $KB^* * \alpha$ is a belief set.
(*2) (Success) $\alpha \in KB^* * \alpha$.
(*3) (Expansion 1) $KB^* * \alpha \subseteq KB^* \# \alpha$.
(*4) (Expansion 2) If $\neg \alpha \notin KB^*$, then $KB^* \# \alpha \subseteq KB^* * \alpha$.
(*5) (Consistency) $KB^* * \alpha$ is inconsistent iff $\vdash \neg \alpha$.
(*6) (Preservation) If $\vdash \alpha \leftrightarrow \beta$, then $KB^* * \alpha = KB^* * \beta$.
(*7) (Conjunction 1) $KB^* * (\alpha \land \beta) \subseteq (KB^* * \alpha) \# \beta$.
(*8) (Conjunction 2) If $\neg \beta \notin KB^* * \alpha$, then, $(KB^* * \alpha) \# \beta \subseteq KB^* * (\alpha \land \beta)$.

The equivalence between our approach and AGM approach is brought out by the following two theorems.

Theorem 6. Let $KB$ a Horn knowledge base with an empty $\mathcal{I}C$ and $+$ be a revision function that satisfies all the rationality postulates ($+1$) to ($+7$). Let a revision operator $*$ on $KB^*$ be defined as: for any sentence $\alpha$, $KB^* * \alpha = (KB + \alpha)^*$. The revision operator $*$, thus defined satisfies all the AGM-postulates for revision ($*1$) to ($*8$).

Proof.

(*1) $KB^* * \alpha$ is a belief set.
This follows immediately, because $(KB + \alpha)^*$ is closed wrt $Cn_P$.
(*2) $\alpha \in KB^* * \alpha$.
This follows from the fact that $+$ satisfies ($+2$).
(*3) $KB^* * \alpha \subseteq KB^* \# \alpha$.

(*4) If $\neg \alpha \notin KB^*$, then $KB^* \# \alpha \subseteq KB^* * \alpha$.
These two postulates follow from ($+4$) and theorem 5.
(*5) $KB^* * \alpha$ is inconsistent iff $\vdash \neg \alpha$.
This follows from from ($+3$) and our assumption that $\mathcal{I}C$ is empty.
(*6) If $\vdash \alpha \leftrightarrow \beta$, then $KB^* * \alpha = KB^* * \beta$.
This corresponds to ($+5$).
(*7) $KB^* * (\alpha \land \beta) \subseteq (KB^* * \alpha) \# \beta$. This follows from ($+6$) and theorem 5.
(*8) If $\neg \beta \notin KB^* * \alpha$, then, $(KB^* * \alpha) \# \beta \subseteq KB^* * (\alpha \land \beta)$.
This follows from ($+7$) and theorem 5. ■
Theorem 7. Let $KB$ a Horn knowledge base with an empty $IC$ and $*$ a revision operator that satisfies all the AGM-postulates (*1) to (*8). Let a revision function $+$ on $KB$ be defined as: for any sentence $\alpha$, $(KB + \alpha)^* = KB^* + \alpha$. The revision function $+$, thus defined, satisfies all the rationality postulates (+1) to (+7).

Proof.
(+1) $P, Ab$ and $IC$ do not change. Obvious.
(+2) $\alpha$ is accepted in $KB + \alpha$. Follows from (*2).
(+3) If $\alpha$ is satisfiable and consistent with $IC$, then $KB + \alpha$ is consistent. Since we have assumed $IC$ to be empty, this directly corresponds to (*5).
(+4) If $\neg \alpha$ is not accepted in $KB$, then $KB + \alpha \equiv KB + \alpha$. Follows from (*3) and (*4).
(+5) If $KB \equiv KB'$ and $\alpha \equiv \beta$, then $KB + \alpha \equiv KB' + \beta$. Since $KB \equiv KB'$ they represent same belief set, i.e. $KB^* = KB'^*$. Now, this postulate follows immediately from (*6).
(+6) $(KB + \alpha) + \beta$ implies $KB + (\alpha \land \beta)$. Corresponds to (*7).
(+7) If $\neg \beta$ is not accepted in $KB + \alpha$, then $KB + (\alpha \land \beta)$ implies $(KB + \alpha) + \beta$. Corresponds to (*8).

4.3 Contraction

AGM puts forward rationality postulates ($-1$) to ($-8$) to be satisfied by a contraction operator on closed set $KB^*$, reproduced below:

$-1$ (Closure) $KB^* - \alpha$ is a belief set.
$-2$ (Inclusion) $KB^* - \alpha \subseteq KB^*$.
$-3$ (Vacuity) If $\alpha \notin KB^*$, then $KB^* - \alpha = KB^*$.
$-4$ (Success) If $\vdash \alpha$, then $\alpha \notin KB^* - \alpha$.
$-5$ (Preservation) If $\vdash \alpha \leftrightarrow \beta$, then $KB^* - \alpha = KB^* - \beta$.
$-6$ (Recovery) $KB^* \subseteq (KB^* - \alpha) + \alpha$.
$-7$ (Conjunction 1) $KB^* - \alpha \land KB^* - \beta \subseteq KB^* - (\alpha \land \beta)$.
$-8$ (Conjunction 2) If $\alpha \notin KB^* - (\alpha \land \beta)$, then $KB^* - (\alpha \land \beta) \subseteq KB^* - \alpha$.

As in the case of revision, the equivalence is brought out by the following theorems. Since contraction is constructed in terms of revision, these theorems are trivial.

Theorem 8. Let $KB$ be a Horn knowledge base with an empty $IC$ and $-$ be a contraction function that satisfies all the rationality postulates ($-1$) to ($-8$). Let a contraction operator $-$ on $KB^*$ be defined as: for any sentence $\alpha$, $KB^* - \alpha = (KB^-\alpha)^*$. The contraction operator $-$, thus defined, satisfies all the AGM-postulates for contraction ($-1$) to ($-8$).
Theorem 9. Let $KB$ be a Horn knowledge base with an empty IC and $-\cdot$ be a contraction operator that satisfies all the AGM- postulates ($-\cdot 1$) to ($-\cdot 8$). Let a contraction function $\dot{-\cdot}$ on $KB$ be defined as: for any sentence $\alpha$, $(KB\dot{-\cdot}\alpha)^* = KB^* - \alpha$. The contraction function $\dot{-\cdot}$, thus defined, satisfies all the rationality postulates ($\dot{-\cdot} 1$) to ($\dot{-\cdot} 8$).

Proof. Follows from Theorem 3 and Theorem 7. ■

5 Realizing Horn knowledge base dynamics using abductive explanations

In this section, we explore how belief dynamics can be realized in practice (see [7,10,11]). Here, we will see how revision can be implemented based on the construction using models of revising sentence and an order among them. The notion of abduction proves to be useful and is explained in the sequel.

Let $\alpha$ be a sentence in $L$. An abductive explanation for $\alpha$ wrt $KB$ is a set of abductive literals $\Delta$ s.t. $\Delta$ consistent with IC and $\Delta \models_P \alpha$ (that is $\alpha \in Cn_P(\Delta)$). Further $\Delta$ is said to be minimal iff no proper subset of $\Delta$ is an abductive explanation for $\alpha$.

The basic idea to implement revision of a Horn knowledge base $KB$ by a sentence $\alpha$, is to realize $Mod(\{\alpha\} \cup IC)$ in terms of abductive explanations for $\alpha$ wrt $KB$. We first provide a useful lemma.

Definition 4. Let $KB$ be a Horn knowledge base, $\alpha$ a sentence, and $\Delta_1$ and $\Delta_2$ be two minimal abductive explanations for $\alpha$ wrt $KB$. Then, the disjunction of $\Delta_1$ and $\Delta_2$, written as $\Delta_1 \lor \Delta_2$, is given as:

$$\Delta_1 \lor \Delta_2 = (\Delta_1 \cap \Delta_2) \cup \{\alpha \lor \beta | \alpha \in \Delta_1 \setminus \Delta_2 \text{ and } \beta \in \Delta_2 \setminus \Delta_1\}.$$ 

Extending this to $\Delta^*$, a set of minimal abductive explanations for $\alpha$ wrt $KB$, $\lor \Delta^*$ is given by the disjunction of all elements of $\Delta^*$.

Lemma 5. Let $KB$ be a Horn knowledge base, $\alpha$ a sentence, and $\Delta_1$ and $\Delta_2$ be two minimal abductive explanations for $\alpha$ wrt $KB$. Then, $Mod(\Delta_1 \lor \Delta_2) = Mod(\Delta_1) \cup Mod(\Delta_2)$.

Proof. First we show that every model of $\Delta_1$ is a model of $\Delta_1 \lor \Delta_2$. Clearly, a model $M$ of $\Delta_1$ satisfies all the sentences in $(\Delta_1 \cap \Delta_2)$. The other sentences in $(\Delta_1 \lor \Delta_2)$ are of the form $\alpha \lor \beta$, where $\alpha$ is from $\Delta_1$ and $\beta$ is from $\Delta_2$. Since $M$ is a model of $\Delta_1$, $\alpha$ is true in $M$, and hence all such sentences are satisfied by $M$. Hence $M$ is a model of $\Delta_1 \lor \Delta_2$ too. Similarly, it can be shown that every model of $\Delta_2$ is a model of $\Delta_1 \lor \Delta_2$ too.

Now, it remains to be shown that every model $M$ of $\Delta_1 \lor \Delta_2$ is either a model of $\Delta_1$ or a model of $\Delta_2$. We will now show that if $M$ is not a model of $\Delta_2$, then

\footnote{An abductive literal is either an abducible $A$ from $Ab$, or its negation $\neg A$.}
it must be a model of $\Delta_1$. Since $M$ satisfies all the sentences in $(\Delta_1 \cap \Delta_2)$, we need only to show that $M$ also satisfies all the sentences in $\Delta_1 \setminus \Delta_2$. For every element $\alpha \in \Delta_1 \setminus \Delta_2$: there exists a subset of $(\Delta_1 \cup \Delta_2)$, $\{\alpha \lor \beta | \beta \in \Delta_2 \setminus \Delta_2\}$. $M$ satisfies all the sentences in this subset. Suppose $M$ does not satisfy $\alpha$, then it must satisfy all $\beta \in \Delta_1 \setminus \Delta_2$. This implies that $M$ is a model of $\Delta_2$, which is a contradictory to our assumption. Hence $M$ must satisfy $\alpha$, and thus a model of $\Delta_1$. ■

As one would expect, all the models of revising sentence $\alpha$ can be realized in terms abductive explanations for $\alpha$, and the relationship is precisely stated below.

**Lemma 6.** Let $KB$ be a Horn knowledge base, $\alpha$ a sentence, and $\Delta^*$ the set of all minimal abductive explanations for $\alpha$ wrt $KB$. Then $\text{Mod}(\{\alpha\} \cup IC) = \text{Mod}(\lor \Delta^*)$.

**Proof.** It can be easily verified that every model $M$ of a minimal abductive explanation is also a model of $\alpha$. Since every minimal abductive explanation satisfies $IC$, $M$ is a model of $\alpha \cup IC$. It remains to be shown that every model $M$ of $\{\alpha\} \cup IC$ is a model of one of the minimal abductive explanations for $\alpha$ wrt $KB$. This can be verified by observing that a minimal abductive explanation for $\alpha$ wrt $KB$ can be obtained from $M$. ■

Thus, we have a way to generate all the models of $\{\alpha\} \cup IC$, and we just need to select a subset of this based on an order that satisfies $(\leq 1)$ to $(\leq 5)$. Suppose we have such an order that satisfies all the required postulates, then this order can be mapped to a particular set of abductive explanations for $\alpha$ wrt $KB$. This is stated precisely in the following theorem. An important implication of this theorem is that there is no need to compute all the abductive explanations for $\alpha$ wrt $KB$. However, it does not say which abductive explanations need to be computed.

**Theorem 10.** Let $KB$ be a Horn knowledge base, and $\leq_{KB}$ be an order among abductive interpretations in $S$ that satisfies all the rationality axioms $(\leq 1)$ to $(\leq 5)$. Then, for every sentence $\alpha$, there exists $\Delta^*$ a set of minimal abductive explanations for $\alpha$ wrt $KB$, s.t. $\text{Min}(\text{Mod}(\{\alpha\} \cup IC), \leq_{KB})$ is a subset of $\text{Mod}(\lor \Delta^*)$, and this does not hold for any proper subset of $\Delta^*$.

**Proof.** From Lemma 6. and Lemma 5., it is clear that $\text{Mod}(\{\alpha\} \cup IC)$ is the union of all the models of all minimal abductive explanations of $\alpha$ wrt $KB$. $\text{Min}$ selects a subset of this, and the theorem follows immediately. ■

The above theorem 10. is still not very useful in realizing revision. We need to have an order among all the interpretations that satisfies all the required axioms, and need to compute all the abductive explanations for $\alpha$ wrt $KB$. The need to compute all abductive explanations arises from the fact that the converse of the above theorem does not hold in general. This scheme requires an universal
order ≤, in the sense that same order can be used for any Horn knowledge base. Otherwise, it would be necessary to specify the new order to be used for further modifying \((KB + α)\). However, even if the order can be worked out, it is not desirable to demand all abductive explanations of \(α\) wrt \(KB\) be computed. So, it is desirable to work out, when the converse of the above theorem is true. The following theorem says that, suppose \(α\) is rejected in \(KB\), then revision of \(KB\) by \(α\) can be worked out in terms of some abductive explanations for \(α\) wrt \(KB\).

**Theorem 11.** Let \(KB\) be a Horn knowledge base, and a revision function \(\vdash\) be defined as: for any sentence \(α\) that is rejected in \(KB\), \(\text{Mod}(KB + α)\) is a non-empty subset of \(\text{Mod}(\lor Δ^*)\), where \(Δ^*\) is a set of all minimal abductive explanations for \(α\) wrt \(KB\). Then, there exists an order \(≤_{KB}\) among abductive interpretations in \(S\), s.t. \(≤_{KB}\) satisfies all the rationality axioms (≤ 1) to (≤ 5) and \(\text{Mod}(KB + α) = \text{Min}(\text{Mod}(\{α\} \cup IC), ≤_{KB})\).

Proof. It is easy to define a pre-order s.t. every model of \(\text{Mod}(KB + α)\) is strictly minimal than all other interpretations. It is easy to verify that such a pre-order satisfies (≤ 1) to (≤ 5). In particular, since \(α\) is rejected in \(KB\), (≤ 3) faithfulness is satisfied, and since non-empty subset of \(\text{Mod}(\lor Δ^*)\) is selected, (≤ 4) is also satisfied. ■

An important corollary of this theorem is that, revision of \(KB\) by \(α\) can be realized just by computing one abductive explanation of \(α\) wrt \(KB\), and is stated below.

**Corollary 1.** Let \(KB\) be a Horn knowledge base, and a revision function \(\vdash\) be defined as: for any sentence \(α\) that is rejected in \(KB\), \(\text{Mod}(KB + α)\) is a non-empty subset of \(\text{Mod}(Δ)\), where \(Δ\) is an abductive explanations for \(α\) wrt \(KB\). Then, there exists an order \(≤_{KB}\) among abductive interpretations in \(S\), s.t. \(≤_{KB}\) satisfies all the rationality axioms (≤ 1) to (≤ 5) and \(\text{Mod}(KB + α) = \text{Min}(\text{Mod}(\{α\} \cup IC), ≤_{KB})\). ■

The precondition that \(α\) is rejected in \(KB\) is not a serious limitation in various applications such as database updates and diagnosis, where close world assumption is employed to infer negative information. For example, in diagnosis it is generally assumed that all components are functioning normally, unless otherwise there is specific information against it. Hence, a Horn knowledge base in diagnosis either accepts or rejects normality of a component, and there is no "don’t know" third state. In other words, in these applications the Horn knowledge base is assumed to be complete. Hence, when such a complete Horn knowledge base is revised by \(α\), either \(α\) is already accepted in \(KB\) or rejected in \(KB\), and so the above scheme works fine.

### 6 Related Works

We begin by recalling previous work on view deletion. Chandrabose [10,11], defines a contraction operator in view deletion with respect to a set of formulae
or sentences using Hansson’s belief change. Similar to our approach, he focused on set of formulae or sentences in Horn knowledge base revision for view update wrt. insertion and deletion and formulae are considered at the same level. Chandrabose proposed different ways to change Horn knowledge base via only database deletion, devising particular postulate which is shown to be necessary and sufficient for such an update process.

Our Horn knowledge base consists of two parts, immutable part and updatable part, but focus is on principle of minimal change. There are more related works on that topic. Eiter is focusing on revision from different perspective - prime implication. Segerberg defined new modeling for belief revision in terms of irrevocability on prioritized revision. Hansson constructed five types of non-prioritized belief revision. Makinson developed dialogue form of revision AGM. Papini defined a new version of Horn knowledge base revision.

We are bridging gap between philosophical work, paying little attention to computational aspects of database work. In such a case, Hansson’s kernel change is related with abductive method. Aliseda’s book on abductive reasoning is one of the motivation keys. Christiansen’s work on dynamics of abductive logic grammars exactly fits our minimal change (insertion and deletion).

In general, our abduction theory is related to Horn knowledge base dynamics (see how abduction theory is related with other applications, respectively, reasoning, update, equivalence and problem solving). More similar to our work is paper presented by Bessant et al. local search-based heuristic technique that empirically proves to be often viable, even in the context of very large propositional applications. Laurent et al. parented updating deductive databases in which every insertion or deletion of a fact can be performed in a deterministic way.

Furthermore, and at a first sight more related to our work, some work has been done on ”core-retainment” (same as our immutable part) in the model of language splitting introduced by Parikh. More recently, Doukari, Özçep and Wu, et al. applied similar ideas for dealing with knowledge base dynamics. These works represent motivation keys for our future work. Second, we are dealing with how to change minimally in the theory of ”principle of minimal change”, but current focus is on finding second best abductive explanation and 2-valued minimal hypothesis for each normal program. Finally, when we presented Horn knowledge base change in abduction framework, we did not talk about compilability and complexity (see the works of Liberatore and Zanuttini).

7 Conclusion

The main contribution of this work lies in showing how abductive framework deals with Horn knowledge base dynamics via belief change operation. We consider the immutable part as defining a new logical system. By a logical system, we mean that it defines its own consequence relation and closure operator. We
presented that relationship of the coherence approach of AGM with this consequence operator is exactly same as our foundational approach, when $IC$ is empty.

We believe that Horn knowledge base dynamics can also be applied to other applications such as view maintenance, diagnosis, and we plan to explore it in further works [9]. Still, a lot of developments are possible, for improving existing operators or for defining new classes of change operators. As immediate extension, question raises: is there any real life application for AGM in 25 year theory? [18]. The revision and update are more challenging in Horn knowledge base dynamic, so we can extend the theory to combine results similar to Konieczny’s [24] and Nayak’s [34].

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