LiKe: Lightweight Certificateless Key Agreement for Secure IoT Communications

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$\label{eq:LiKe} LiKe ~{\rm is} \\ {\rm a~lightweight~pairing-free^1~certificateless~key~agreement~protocol} \\$

¹Pairing-based cryptography: pairing between elements of two cryptographic groups to a third group with a mapping $e: G_1 \times G_2 \to G_T$.

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- ephemeral cryptographic materials
- support for intermittent connectivity with the Trusted Third Party(TTP)
- lightweight *re-keying* operations
- robustness against impersonation attacks (even when information stored on the TTP are leaked)
- formally secure (via the *ProVerif* tool)
- compatible with real IEEE 802.15.4 devices
- suitable for integration in Zigbee 3.0
- bandwidth and energy efficient

- Certification Management in traditional public key infrastructure (PKI) is inefficient
- Key-escrow problem in Identity Based Encryption

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Certificateless Cryptography

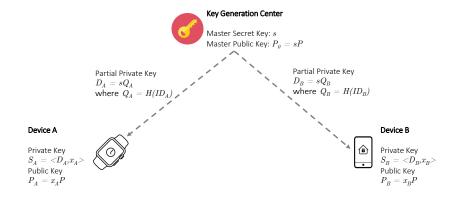


Figure: Certificateless Cryptography - An overview.

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Scenario

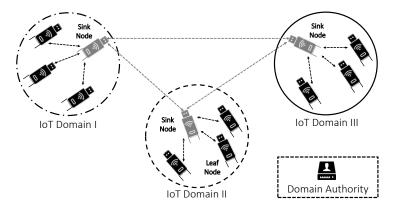


Figure: Reference Scenario: multiple IoT domains, managed by a unique Domain Authority (DA), where each domain is organized in a single sink IoT node and several leaf IoT nodes.

Pre-Deployment Phase

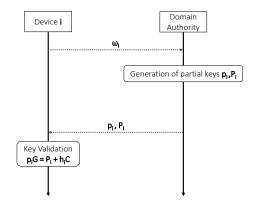


Figure: LiKe: Sequence Diagram of the Setup Phase.

$$\begin{aligned} \omega_i &= (ID_i \parallel t_i \parallel X_i) & P_i &= r_i G \\ h_i &= H(\omega_i \parallel P_i) & p_i &= r_i + h_i c \mod r \end{aligned}$$

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Key Agreement Protocol

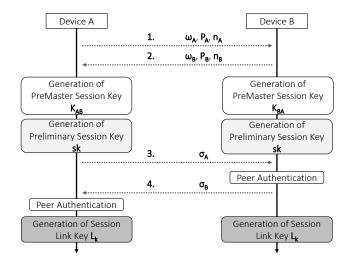


Figure: LiKe: Sequence Diagram of the Key Agreement Phase.

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Protection Against Leakage of Secret DA Information

Ephemeral Cryptography Materials

Protection Against Man-in-the-Middle Attacks

Protection Against Replay Attacks

Protection Against Known-Key Attacks

```
:~$ proverif LiKe_test 1.pv | grep "RESULT" | nl
1 RESULT not attacker(XA[]) is true.
2 RESULT not attacker(XA[]) is true.
3 RESULT not attacker(pA[]) is true.
4 RESULT not attacker(pB[]) is true.
5 RESULT inj-event(end_LiKe_B(x)) ==> inj-event(begin_LiKe_A(u,v_13,y,w,n)) is true.
6 RESULT inj-event(end LiKe A(x 14)) ==> inj-event(begin_LiKe_B(u) 15,v 16,v 17,w 18,n 19)) is true.
```

(a) Test 1: Security Verification of LiKe when the private keys of the involved devices are assumed to be secret.

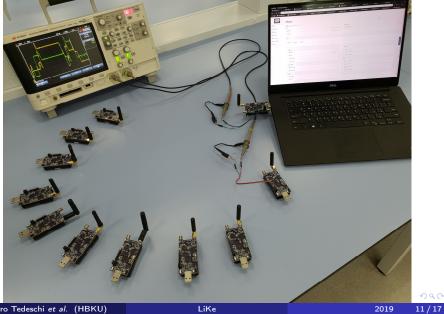
```
:-$ proverif Like_test_2.pv | grep "RESULT" | nl
1 RESULT not attacker(XA[]) is true.
2 RESULT not attacker(XA[]) is true.
3 RESULT not attacker(pA[]) is false.
4 RESULT not attacker(pB[]) is false.
5 RESULT inj-event(end_Like_B(X)) ==> inj-event(begin_LiKe_A(u,v_13,y,w,n)) is true.
6 RESULT inj-event(end_LiKe_A(X_14)) ==> inj-event(begin_LiKe_B(u_15,v_16,y_17,w_18,n_19)) is true.
```

(b) Test 2: Security Verification of LiKe when the partial private keys of the involved devices maintained by the DA are assumed to be leaked to the adversary.

Figure: Output provided by ProVerif.

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Experimental Testbed



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- LiKe Protocol Duration: 340.478ms
- Preliminary Session Key Generation: 243.392ms
- Mutual Authentication and Session Key Generation: 97.086ms

[Comparison] - MAC-Layer Exchanged Messages

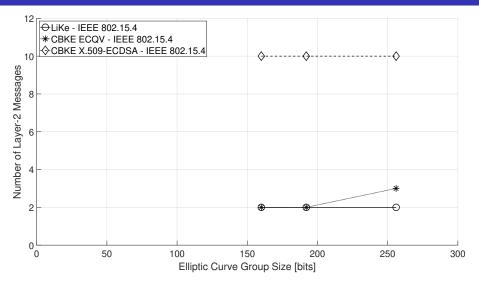


Figure: MAC-layer messages required to complete the key agreement, considering X.509-ECDSA, ECQV, and LiKe.

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Feature	X.509-ECDSA	ECQV	LiKe	
Robustness to Man-In-The-Middle	1	1	1	
No. of messages to detect an attack	9	2	2	
Message Overhead per Key Agreement Instance	10	3	2	
Energy Consumption per Key Agreement Instance (mJ)	38,953	36,080	35,726	
Robustness to DA Secret Information Disclosure	×	×	1	

Figure: Comparison of LiKe against CBKE Approaches based on X.509-ECDSA and ECQV, assuming a 256-bit Elliptic Curve.

Energy Consumption

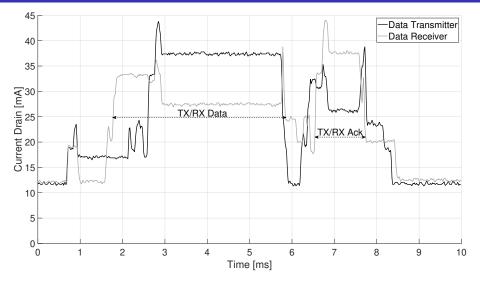


Figure: Current Consumption of a Data Transmission and a Data Reception Operation within the duration of a IEEE 802.15.4 time-slot (10 ms). 2019 15 / 17

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LiKe

Feature	[17]	[18]	[19]	[20]	[21]	[22]	[23]	[24]	LiKe
Pairing-free	X	X	1	X	1	1	X	1	1
Non-persistent Connection with Domain Authority	X	X	1	X	X	~	1	1	1
Ephemeral Cryptograhy Material	X	X	1	X	1	1	X	X	1
Integration in a Real IoT Enabling Technology	X	X	X	X	X	~	X	X	1
Implementation on Real IoT Devices	X	X	X	X	X	1	X	X	1
Real Performance Evaluation	X	X	X	X	X	~	X	X	1
Energy Friendly Approach	-	-	-	-	-	X	-	-	1
Suitability for Autonomous IoT Domains	X	X	X	X	X	X	X	X	1
Lightweight Re-Keying	X	X	X	X	X	X	X	X	1
Open-Source Code	X	X	X	X	X	X	X	X	1

Figure: Comparison of LiKe against state-of-the-art approaches using **CL-PKC** techniques. A \checkmark symbol indicates the fulfillment of a particular feature, a \varkappa symbol denotes the miss of the feature, while the symbol — indicates that the feature is not applicable.

Any Questions? Thank you!



Image: A matrix

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