Maude Object-Oriented Action Tool

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Introduction

- AS is a formal framework.
- Its notation presents good reusability and extensibility properties.
- However, AS lacks of syntactic support for the definition of modules and libraries.
- A modular approach introduces modules into AS descriptions.
- Based on the modular approach and object-orientation, OOAS was proposed.
Introduction

- OOAS was defined using SOS transitions.
- No tool for OOAS had been implemented yet.
- Here we present MOOAT:
  - the first OOAS implementation;
  - changes on OOAS syntax;
  - OOAS implementation using MSOS;
  - how the tool has been built using Full Maude and Maude MSOS Tool.
- In addition, we show how to use Mosses’ *constructive* approach with OOAS as a case study of MOOAT.
Maude MSOS Tool

- It is an executable environment for describing MSOS specifications.
- It is a conservative extension of Full Maude that compiles MSOS specifications into Rewriting Logic.
- Uses an adopted MSDF syntax:
  - Modular SOS Definition Formalism - Mosses.
- MMT is quite useful:
  - the abstract syntax is defined using BNF;
  - the semantics is given by *labeled transitions*, which contains the necessary semantic components;
  - a specification might be organized by modules.
MMT limitations

- **pre-regularity check**
  - A term must contain just one sort.

  Wrong: Storable ::= Int . Value ::= Int .

  Correct: Storable ::= Int . Value ::= Storable .

- **ad-hoc overloading**
  - If the sorts in the arities of two operators with the same syntactic form are pairwise in the same connected components, then the sorts in the coarities must likewise be in the same connected component.

  Wrong: Exp ::= sum (Exp, Exp) | Int | Id . Env = (Id, Int) Map . Sto ::= (Int, Int) Map .

  Correct: Exp ::= sum (Exp, Exp) | < Int > | Id . Env = (Id, Int) Map . Sto ::= (Int, Int) Map .
Object-Oriented Action Semantics

- Method to organize AS specifications.
- Goals:
  - enhance extensibility and reusability;
  - split semantic descriptions into classes;
  - treat semantic functions as methods.
- OOAS extends AN:
  - to accept classes and operators based on object-orientation;
  - a class is split into two parts: syntax and semantics;
  - an OOAS description is a hierarchy of classes.
Class Command
syntax:
   Cmd
semantics:
   execute : Cmd → Action
End Class

Class While
extending Command
using E:Expression, C:Command
syntax:
   Cmd ::= "while" E "do" C
semantics:
   execute [[ "while" E "do" C ]] =
      unfolding
         evaluate E then
         execute C and then unfold else complete
End Class
State class

- It is the base-class for all OOAS classes.
- Since all classes are under a pre-defined root class, they belong to a hierarchy of classes.
- It is responsible to implement the attributes and operations that are used in those classes specifications.
- OOAS actions are exactly the same actions of AS.
- **State** attributes:
  - `transients`;
  - `bindings`;
  - `storage`.

Some operations are available to handle State attributes. They are classified according to AS facets:

- basic - control flow;
- functional - transients;
- declarative - bindings;
- imperative - storage;
- reflective - abstractions;
- hybrid - actions that deal with more than one facet.

No operation was defined to deal with communicative facet.

The formal semantics of each operation has been defined using SOS.
Constructive Action Semantics

- Constructs are defined separately and they can be used in several specifications by deriving constructs.
- The set of constructs is not minimal, in a specification just the necessary constructs are included in it.
- In this way, all available constructs do not need to be included in every definition.
- The constructive approach provides a huge flexibility on the definition of programming languages.
  - In this regard, it has been used with Modular Action Semantics and Modular Structural Operational Semantics.
Constructs

- A construct is the formal specification of a certain programming language feature.

- Abstract and Concrete Constructs:
  - while (Exp) do Cmd (concrete);
  - cond-loop(Exp, Cmd) (abstract);
  - a language-independent prefix notation is used to define abstract constructs.

- Basic and Derived Constructs:
  - basic constructs represent common features that have the same interpretation and are found in many programming languages;
  - derived constructs represent particular features regarding the programming language;
  - derived constructs are obtained by combining abstract constructs.
The CAS goal is to map concrete constructs from programming languages to a set of basic abstract constructs.
Maude Object-Oriented Action Tool

- **MOOAT**
  - The first executable environment for describing programming languages using OOAS.
  - MMT has been used to specify the Action Notation.
  - Full Maude has been used to implement the Classes Notation.
  - MOOAT development was inspired by MAT (Maude Action Tool).

- MOOAT is a conservative extension of Full Maude and MMT
  - It has the same limitations as Full Maude and MMT.
  - Regarding these limitations, its syntax has some differences if compared to the formalism one.
  - Yet, MOOAT syntax and OOAS syntax are quite similar.
An OOAS specification is a finite set of classes. The classes create the necessary hierarchy. Syntax trees in MMT BNF style are used. The semantics is given by Action Notation.
Data Notation

DN has been implemented using 1 MSDF module and 1 Maude system module.

- **MSDF - DataNotation**
  - The sorts of data used by AN are defined.
  - Action, Yielder and Data - sorts implemented by MOOAT.
  - Int and Boolean - sorts implemented by MMT and used by MOOAT.

- **Maude - DATANOTATION**
  - Implement the function $\text{op}_\langle \_ :<\_ : \text{Data DataSort} \to \text{Bool}.$
  - Used to verify if a datum belongs to a sort.
  - $<1>$ :< value returns true.
  - $<$ cell (1) > :< abstraction returns false.

Some auxiliary operations has been defined:

- $\text{sum} \ (\text{Yielder}, \text{Yielder})$. 
Action Notation

- AN has been split into MSDF modules:
  - the modules have been used to implement the available facets;
  - AN specifies the indexes used in the transition labels.

(msos BasicSyntax is
see BasicData.

```
Action ::= Action or Action | fail | commit |
          Action and Action | complete |
          indivisibly Action |
          Action and 'then Action |
          Action trap Action | escape |
          unfolding Action | unfold | diverge.
```

Yielder ::= the DataSort yielded by Yielder |
          nothing | Data.

sosm)
Action Notation

(msos FunctionalOutcomes is
  see BasicOutcomes.
  see FunctionalData.
Completed.
Terminated ::= Completed.
Completed ::= completed.
Completed ::= gave (Data).
gave (none) : Action --> completed.
sosm)

(msos BasicConfigurations is
  see BasicSyntax.
  see BasicOutcomes.
Action ::= Terminated |
  Action @ Action.
sosm)

(msos FunctionalLabels is
  see BasicLabels.
  see FunctionalData.
Label = {data : Transients,
  data’ : Transients, ...}.
sosm)
Each facet implements the necessary transitions.

We have basically 3 kinds of transitions:

- those transitions that neither change or use labels’ components:
  
  \[
  \text{Action1} \xrightarrow{\ldots} \text{Action’1} \\
  \text{Action1 and Action2} : \text{Action} \xrightarrow{\ldots} \text{Action’1 and Action2} .
  \]

- those transitions that just use labels’ components:
  
  \[
  \text{regive} : \text{Action} \xrightarrow{\text{data = Transients, data’ = Transients,}} \text{gave (< Transients >)} .
  \]

- those transitions that either change and use label’s components:
  
  \[
  \text{gave (Data1) and gave (Data2)} : \text{Action} \xrightarrow{\text{data = Transients, data’ = Transients’’,}} \text{gave (concatenation(Data1, Data2))} .
  \]
State class

- It is represented by a Maude system module called STATE.
  - Since MOOAT classes are translated to system modules and also because all MOOAT classes are subclasses of STATE as well as in OOAS.
- And also by Classes Notation.
- MOOAT State class has 5 attributes.

(mod STATE is
  including ActionNotation.
  op init-rec : -> Record.
  eq init-rec = {commitment = false,
                 unfolding = fail,
                 data = (void).Map{Int,Datum},
                 bindings = (void).Map{Token,Data},
                 storage = (void).Map{Cell,Data}}.

  var A : Action.
  var S : Storage.
  eq new-cell (S) = cell (length (S) + 1).
  op compute : Action -> Conf.
  eq compute (A) = < A ::: 'Action, init-rec >.
endm)
Classes Notation

- CN has been developed similar to:
  - Full Maude object-oriented modules;
  - MMT MSDF modules.

- When CN syntax is used its constructs are translated to code that Maude is able to interpret.

- MOOAT classes are composed basically by 3 parts:
  - extends part;
  - syntactic part;
  - semantics part.

- The use of these parts might be optional or sequential. However, a class must has at least one of these 3 parts.
Classes Notation

- Each class part is translated as follows:
  - subclasses definitions are simply translated to lines that use Maude’s directive including;
  - syntactic definitions use the BNF style introduced by MMT - trees and syntactic sorts are translated to Maude’s operations, sorts and subsorts;
  - in the semantic definitions, methods and objects are translated respectively to Maude’s equation sets and meta-variables.

- A methods environment is created by the inclusion of the converted MOOAT classes into Maude’s module database.

```maude
rl < O : X@Database | db : DB, input : ('class_is_endclass[T, T'], step-flag : B,
output : nil, default : MN, Atts > =
< O : X@Database | db : mooat-proc-unit('class_is_endclass[T, T'], step-flag(B), DB),
input : nilTermList, step-flag : B, output : ('Introduced 'OOAS 'class
header2QidList(parseHeader(T)) '
), default : parseHeader(T), Atts > .
```
Constructive Object-Oriented Action Semantics

- COOAS = CAS + OOAS.
- Case study of MOOAT.

(class Cmd is
  Cmd .
  absmethod execute Cmd -> Action .
endclass)

(class Cmd/While is
  extends Cmd .
  Cmd ::= cmd-while (Exp, Cmd) .
  method execute (cmd-while(E:Exp, C:Cmd)) = unfolding
    (evaluate E:Exp then
     ((execute C:Cmd and then unfold) else complete)) .
endclass)
μ-Pascal is a toy language fairly similar to the imperative language Pascal.

It contains basically commands and expressions.

It is being described just by extending the necessary COOAS classes to compose the language.

(class Micro-Pascal is
  extends Exp/Val, Exp/Val-Id .
  extends Exp/Sum, Exp/Sub, Exp/Prod .
  extends Cmd/Assignment, Cmd/Repeat, Cmd/While .
  extends Cmd/Sequence, Cmd/Cond .
  extends Dec/Variable, Dec/DecSeq .
  extends Prog .
endclass)
We have presented MOOAT, the first OOAS implementation.

MOOAT is a conservative extension of Full Maude and MMT.

The previous OOAS specification using SOS has been rewritten using MSOS and implemented using MMT.

A language to specify OOAS classes has been built in Maude.

MOOAT is the update of OOAS since its formal specification has been translated from SOS to MSOS.

MOOAT notation covers a rich set of actions and action combinators, which includes the complete OOAS notation.
COOAS represents a novel view to the OOAS system. This is the first time that CAS and OOAS are combined. The modularity aspects observed in OOAS has been improved due the introduction of constructs into it. Since such combination is capable to describe syntax-independent specifications of programming languages. As future work:
- we would implement the communicative facet of Action Notation;
- we would trace a careful comparison between COOAS and a library of OOAS classes called LFLv2.
Thank you

- MOOAT can be found at:
  - http://www.inf.ufpr.br/murbach/mooat/
- Please, send any questions or comments to:
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