Admission Control for Distributed Complex Responsive Systems

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Abstract

The software in modern systems has become too complex to make accurate predictions about their performance under different configurations. Real-time or even responsiveness requirements cannot be met because it is not possible to perform admission control for new or changing tasks if we cannot tell how their execution affects the other tasks already running. Previously, we proposed a resource-allocation middleware that manages the execution of tasks in a complex distributed system with real-time requirements. The middleware behavior can be modeled depending on the configuration of the tasks running, so that the performance of any given configuration can be calculated. This makes it possible to have admission control in such a system, but the model requires knowledge of run-time parameters. We propose the utilization of machine-learning algorithms to obtain the model parameters, and be able to predict the system performance under any configuration, so that we can provide a full admission control mechanism for complex software systems.

1. Introduction

Traditionally, real-time systems have mainly been used in tightly controlled environments, such as mission-critical and embedded systems. These systems are usually characterized by a set of well-known tasks that have to be executed by a scheduler. The tasks are modeled after the time it takes to execute them and the period with which they are to be executed (or, more generally, a deadline by which the task must have finished execution). The scheduler follows a policy that establishes which task has to be executed next, such that all tasks are finished before their corresponding deadlines. Examples of such policies are Rate Monotonic Analysis (RMA) or Earliest Deadline First (EDF) [1]. In the embedded system case, the tasks to execute are known a priori, having the system fully specified and working by design. In scenarios where more flexibility is allowed, the arrival of tasks for execution is not predefined, although their execution characteristics are specified. In that case a mechanism for admission control is needed, such that the scheduler may refuse to execute a task if sufficient resources are not available. This lets the tasks already scheduled be processed knowing that their requirements will be met. Given the scheduling policy and the specification of the task execution characteristics, a calculation can be carried out on whether a new task can be scheduled on the available resources. This calculation is at the core of any admission control mechanism.

Recently, the interest has grown for large distributed enterprise systems to become more “real-time”, i.e., the focus has moved from pure average performance to a more controlled evaluation of the performance, where the worst case matters as much as the average. However, software running on these systems cannot be fully specified in contrast to the embedded counterpart: it is difficult to know at all times what is actually running on a host (background tasks and OS daemons may execute at unpredictable times); even when the applications are known, their characteristics are difficult to describe; traditional techniques for real-time analysis do not fit well with the very complex software stacks run today.

Our attempt to make the problem more tractable has been the introduction of real-time messaging middleware [2] to build distributed responsive systems. If all or most applications running in a system can conform to a common API and programming model, the execution of the applications can be largely under the control of the middleware. Given that, the middleware has the possibility to play the role of the scheduler in the real-time system, with the applications being the tasks. The middleware can account for the temporal consumption of resources by the applications, enforcing a certain partition of the resources much like the scheduler in a
real-time system. Thus, provided that the middleware manages the execution of all (or the most important) activities in the system, and given a mathematical model of the middleware execution, the performance of the system can be calculated in a given configuration. This allows us to have admission control in a complex real-time system, because we are able to determine whether the introduction of a new application violates the performance requirements of the other tasks already running. Still, we found in previous work [3] that, even if the middleware allows the modeling of the system, the complexity of the underlying technology (e.g., garbage collection) makes it impossible to have a closed-form mathematical model of the system that is valid in all circumstances. We could, however, obtain a family of models that depend on some parameter that is determined by the configuration of the system. In this paper, we build up on that modeling and propose a new admission control mechanism based on the estimation of the parameters of the model through machine-learning techniques. For simplicity, in this paper we consider only one parameter.

As we show below, the relation between the configuration of the system and the values of the parameter defining the model is highly nonlinear. The values of the parameter depend on complex interactions at all levels of the software stack and between the nodes. It is not evident to establish a priori the parameter value corresponding to a given configuration, except for the simplest configurations. However, we assume that the configuration and the parameter describing the performance of the system cannot be randomly related, but are subject to some complex logic. We attempt to grasp this logic by use of machine-learning (ML) techniques. We decide to model our problem as a classification system, using decision trees. The objects for classification are the configurations of the system, and the resulting classes are ranges of values for the parameter of interest in the model.

There are a number of ML methods that we could have used, such as neural networks or clustering algorithms. We decided on decision trees for the following reasons: Firstly, there is extensive literature on satisfactory results in a number of fields where decision trees have been applied to complex nonlinear problems. Secondly, the extensive use of decision trees has provided mature algorithms that offer very good performance, a critical feature in a responsive or time-critical system, where decisions need to be made fast enough. Finally, the output of the decision tree creation algorithm is a human-readable structure. This is an invaluable analysis tool for a system administrator, who can obtain insight into the behavior of the complex system by looking at the decisive attributes in the decision tree.

Section 2 offers an overview of related technologies in real-time complex systems, ML algorithms and their applications. Section 3 describes the basics of decision tree algorithms. Section 4 introduces the modeling of the real-time publish/subscribe middleware and the need for a complex nonlinear problem solver. Section 5 describes the whole system design and how it works. Section 6 shows the accuracy of the modeling used for the admission control mechanism and its suitability for responsive systems.

2. Related Work

Quite some work has been developed on admission control for QoS-enabled networking. There are obvious similarities between the purpose of QoS networking and that of our middleware: some resources (bandwidth resp. CPU) are to be shared among competing entities (flows resp. tasks), and the resources have to be allocated according to some reservation (expressing the need for resources). An admission control algorithm for real-time traffic that takes advantage of scheduling information is presented in [4]. Linear programming is used to obtain an admission control policy that preserves the guaranteed QoS and maximizes utility. The authors of [5] measure the maximal rate envelopes of flows. With those measurements, they can predict with some measurable certainty the schedulability of new flows, which are admitted provided the QoS of all flows is respected. Reference [6] describes an admission control algorithm for delay-bounded packet networks that provides a predictive service.

The authors of [7] propose real-time middleware to facilitate the use of the real-time QoS capabilities of the OS by the applications. Thus their approach is different in that they assume the presence of an OS with real-time capabilities, for which the middleware is merely an enabler for applications.

The real-time literature for complex systems can for the most part be described as an extension to distributed systems of the task-scheduling problem [8]. The assumptions about the tasks are the same as in embedded systems: well-known execution time and well-known available resources. In the middleware literature, QoS requirements are generally used as matching criteria for two end-points to connect [9], but without specifying the mechanisms to actually enforce the QoS requirements. These systems must assume over-provisioning of resources, a scenario where admission control is clearly not required. Other approaches more concerned with actual system performance are adaptive
bandwidth reservation schedulers [10]: Reservation-based schedulers reserve a fraction of the total CPU to each task, but they assume that the bandwidth given to a task is sufficient to fulfill its timing constraints, so that no admission control is needed. Instead of preemptively controlling access of new tasks, the authors propose to use adaptive reservation for the tasks, such that each allocated fraction of the CPU is determined by a closed control loop. The authors of [11] also present an on-line approach to adapt the QoS of distributed real-time systems, and use distances between states to define the optimal operational point. Control theory is used to bring the system to a desired point of operation, while maintaining some system performance level. The system attempts to allocate resources for all tasks according to their needs, but it can not be ensured that the requirements of all tasks running can be satisfied. In such systems, even if a task is terminated because its requirements cannot be met, some tasks may in the meantime have been negatively affected by the resource allocation.

The seminal work of Quinlan explaining the ID3 algorithm [12] is a recommended reading for an in-depth understanding of decision trees. The algorithm we use in this paper, C4.5 [13], is one of the most widely used.

3. Decision Trees

Decision trees are general-purpose knowledge-based systems that classify objects into a set of classes according to the attributes of those objects. The decision tree itself represents the knowledge acquired by the system. The leaves of the tree correspond to classes. Every other node in the tree is a test on a given attribute of an object, and there is a branch coming out of a node for every possible outcome of the test. A branch is directed to a leaf, classifying the object, or to another node, where the classification process continues.

Decision trees are useful for expressing nonlinear relations between a set of objects and a set of classes through the attributes of the objects. The usual procedure is that of having a pool of observed data, composed of a set of objects with given attributes and the objects’ classification, which is observed and not systematically obtained. The intention is then to be able to classify another object, given its attributes, following the same logic that permeates the observed data. In other words, the aim is to accurately predict the class corresponding to an unobserved object.

The task of building decision trees from already classified objects and their attributes is in general NP-Complete. There are a number of greedy algorithms (such as ID3 or C4.5) that can efficiently compute good approximations to the optimal solution. The main idea is to create the simplest decision tree that correctly classifies all the observed data. This is achieved by making the branching decision at each node based on the attribute that reduces the most the entropy of the subset of objects classified under that node so far. The entropy of a set of objects increases with the number of classes the objects of the set belong to. When a node produces 0 entropy for all objects classified under that node, it is said to have reached purity, as all those objects correspond to the same class. That node is a leaf node for the decision tree. A detailed explanation of the general principle and the ID3 algorithm in particular can be found in [12]. For our experiments, however, we use C4.5, which is based on ID3. C4.5 solves a number of problems present in ID3: C4.5 avoids over-fitting of the tree to the observed data by pruning the tree; C4.5 allows the description of attributes by a range of real numbers instead of just discrete values, a property that makes our approach feasible. For our experiments we have used the C4.5 algorithm implementation by Quinlan [13], which is available from [14]. In Section 6 we study its performance.

4. Modeling Complex Middleware

Current software stacks are built from heterogeneous components whose internals are usually not fully understood, and the interactions in a distributed system are not always predictable. This makes the extraction of a useful model that provides any insight into the behavior of the system virtually impossible. This in turn prevents predicting how the system will be affected by the addition of a new task, making it impossible to decide whether the task should be accepted. This is why we propose a middleware that takes care of the execution of the different parts of the system and provides a tractable model that can be analyzed, such that admission control can be implemented. A description of our middleware can be found in [2]. At the core there is a scheduler that assigns resources to tasks (details can be found in [3]), described here for completeness. This paper focuses however on the admission control mechanism. In a nutshell, in any node of the system, credits are assigned to tasks, which implement the application logic and communication means. The total credits allocated to a task are a fraction of a time period $T$. The CPU share is expressed as the execution time every $T$. The scheduler always chooses to schedule the task with the highest number of credits that has a nonempty event queue, i.e., that has work to do. When a task is scheduled,
the run time the associated handling function needs is measured and deducted from the total number of credits for that task. The credits of all tasks are reset when all tasks with work to do have no credits left.

This scheduler enables a processor time allocation that follows a Weighted Max-Min policy, which can be expressed using the following mathematical model: On an N processor machine assume that the share allocated to task \( i \) is \( 0 \leq z_i < 1/N \) (tasks are processed by a single scheduler thread) and that \( s_i \) is the number of events arriving at the task queue per second. The number of events \( r_i \) serviced for a task \( i \) can be viewed as containing two components: the rate guaranteed by the share and any spare capacity that a task can take advantage of. Assuming that the cost to process an event is constant, we denote the total maximum achievable rate by \( R^{\text{max}} \) events per second (evt/s). \( R^{\text{max}} \) is a function of system settings and the parameter the model depends on. For a given setting of \( z_i \), \( s_i \) and \( R^{\text{max}} \), the expected rate \( r_i \) is:

\[
r_i = \text{Min} \left( s_i, \frac{R^{\text{max}}}{\frac{z_i}{z_i} - \sum_{j \neq i} s_j} \right)
\]

(1)

where \( u_i = \frac{s_i}{\sum_{j \neq i} s_j} \), \( x_i = \frac{u_i}{z_i} \) and we define \( u_i \) to be the fraction of resources that a topic is attempting to consume and \( x_i \) to be the ratio of \( u_i \) to \( z_i \). Eq. (1) states that the share that a task gets is its adjusted share of what is left over after tasks with higher \( x_i \) have been allocated. The amount of additional unused capacity available to a task is the capacity left over by tasks that are less demanding (lower \( x_i \)) than the task itself.

In order to check the fidelity of the model to the real implementation of the scheduler, we compare the throughput \( o_i \) obtained on each of the \( t \) tasks with the corresponding throughput calculated by the model. We quantify the difference between the \( t \)-dimensional vector of event rates \( \bar{r} = (r_1, r_2, ..., r_t) \) predicted by the model and the vector of actual rates \( \bar{o} = (o_1, o_2, ..., o_t) \) obtained in the experiment by calculating the Euclidean distance between these two vectors, \( e = ||\bar{r} - \bar{o}|| \). The value \( e \) by itself is not a good indicator of the performance of the model, because it depends on the magnitude of \( \bar{r} \) and \( \bar{o} \). To provide some meaningful reference, we compared this distance to the average distance between two randomly chosen points in the same vector space. We perform a number of experiments that cover a range of system configurations. We vary the following parameters: the sum of the throughput being processed by all tasks, which is important with respect to \( R^{\text{max}} \); the number of tasks, which increases the unaccounted overhead; the task shares, and the event rate produced by the tasks, which is in principle independent of the share allocated to that task. We produce 50 different, randomly generated task rate and share settings for each total throughput (30,000, 100,000 and 200,000 evt/sec) and number of tasks (10, 20, 50 and 100), providing a total of 600 scenarios under test using two scheduling threads on a two-core machine. We obtain the \( R^{\text{max}} \) parameter value from each experiment.

The results in Fig. 1 show the cumulative distribution function (CDF) of the relative error when estimating the throughput in the experiments. The model reflects the behavior of the actual system with a very high degree of accuracy. The larger the number of tasks, the lower the fidelity. But even for 100 tasks, the prediction error in the worst case is only 10% of the random predictor error reference. For 10 tasks, almost 90% of the settings’ behavior is described 1,000 times more accurately by our model.

5. System Description and Utilization

Overview. Our middleware provides an execution environment for complex distributed time-sensitive applications. These applications have some timing and performance requirements, which have to be met by using the amount of resources reserved through the middleware.

Our utilization model is such that each application (a task) requests a share of the resources, intended to process a certain rate of events per second. The requirements of a set of tasks form a configuration of the middleware. A configuration is invalid if the amount of resources requested is greater than the total available. Assume a running system in some

![CDF (2 threads) of the prediction error for 100Kevt/sec](image-url)
valid configuration. A new task is introduced (or the requirements of an existing one are changed), resulting in a new configuration \((configuration_x, \text{param})\) in Fig. 2). As can be seen in Fig. 2, the output of the admission control system is either to accept or refuse the new configuration. The admission control mechanism has to decide whether the new configuration can be applied to the system, given the resulting performance. To do so, the system needs to 1) determine the performance resulting from the application of the configuration and 2) decide whether that performance is satisfactory. The model (\(Middleware_{model}\) in Fig. 2) needs a run-time parameter that depends on the configuration. The value of this parameter \((\text{param\_value}_x, \text{param})\) in Fig. 2) is obtained from the new configuration using a knowledge system, in our case a decision tree. Then a prediction for the performance of the system under the new configuration \((model\_\text{performance\_result})\) can be obtained.

Defining whether the performance of the system for a given configuration is satisfactory may be performed in a number of ways: in the most strict case, the performance is only acceptable when all the individual tasks in the system meet their performance requirements. In a more flexible solution, a system administrator may decide that some tasks are more important than others, and thus the configuration may be changed to meet the requirements of some tasks, even at the expense of having some other tasks penalized. We assume to have these criteria available.

![Figure 2. Diagram of the admission control system using the real-time middleware.](image)

The knowledge needed to estimate the run-time parameter of the model is obtained using an ML algorithm. The introduction of ML in the system design means that there has to be a training phase in which the ML algorithm learns which are the criteria linking one configuration to a given value of the model parameter. This training phase may be run off-line, by having the system actually run in a number of configurations and measuring the corresponding values of the parameter of interest (top of Fig. 2). Then this data in the form of \((configuration_x, \text{param\_value}_x, \text{param})\) pairs is fed to the ML algorithm (C4.5 in our case) to obtain the knowledge system (decision tree) used in the admission control mechanism. Note that the data set may be arbitrarily large, but the resulting decision tree does not grow proportionally, because it is in effect a summary of the information contained in the data set. We argue that running the training phase is not an exceptionally constraining requirement, as any software platform needs to undergo a test and verification phase before it is deployed in production; even more so if we speak of software with subtle performance requirements.

It would also be possible to have the decision tree recalculated periodically as new information about the performance of the system is collected. We evaluate the feasibility of this approach in Section 6, by measuring the time required to build a tree.

**System Calibration (Learning Phase).** The calibration consists of gathering information about the behavior of the system under a variety of configurations. Assume that there is a set of configurations of interest such that it represents the expected range of operation of the system. When the system is being calibrated, under each configuration, the behavior of the system is captured by registering the value of the parameter of interest for the middleware model. In this way, at the end we obtain a set of configuration and parameter value pairs. This data is fed to the C4.5 ML algorithm, which creates a decision tree that reflects the nonlinear relation between the configuration and the model parameter value. The configurations are thus the objects to be classified by the decision tree, while the parameter values are the classes. The model parameter value of a configuration is obtained by classifying the configuration with the decision tree.

A configuration is defined as follows: On an \(N\) processor machine, assume that the share allocated to task \(i\) is \(0 \leq z_i < 1/N\) and that \(s_i\) is the number of events arriving at the task queue per second. The configuration for a system with \(t\) tasks is then given by a vector \(W = [(z_0, s_0), \ldots, (z_{t-1}, s_{t-1})]\). Note however that, given that we assume the cost of processing an event to be equal for all tasks, there are different configurations that are equivalent. In fact, as far as the middleware is concerned, all permutations of \(W\) are exactly the same. We therefore need to represent all permutations of \(W\) as a single object, such that the classifier does not become unnecessary.
ily complex. In this way, relations between similar objects (configurations) are captured by the common representation. Thus, the common object representation of all permutations of \( W \) is defined by a vector of pairs: \( V = [(z_0^p, s_0^p) \ldots (z_j^{p-1}, s_j^{p-1})] \), such that \( \forall p, q, p < q \rightarrow s_p^q > s_q^p \lor s_p^q = s_q^p \lor s_p^q > s_q^p \). This means that the first setting in the configuration vector is the one having the largest event rate, the second has the second largest rate, and so on. If two settings have the same event rate, the order is decided by having first the one with the larger share. Each configuration is an object with \( 2t \) attributes, to be classified by the decision tree.

The run-time parameter of the model can have any value within a range in a continuous space, but decision trees have discrete classifications as output. To overcome this apparent mismatch, we divide the range of the parameter value into categories of equal size. The range can be estimated from the calibration phase itself, as the various configurations under test, if properly chosen, provide a good indication of the values expected for the parameter during normal execution. This range may be extended to account for outlier values not seen during calibration. In effect, we are quantizing the estimation of the parameter value. The less quantization error we want to introduce, the smaller the classes and thus the more classes we have to define in the output of the decision tree. For simplicity, we opt for a uniform quantization scheme, although with some models a nonuniform scheme may yield better results.

To summarize, during the calibration period a certain number of configurations and their corresponding \( R_{\text{max}} \) are obtained. We denote the \( k \)th normalized configuration as vector \( V_k \). Hence the result of the calibration period is a set \( C \) of pairs of the form \( (V_k, R_{\text{max}}^k) \), which are used to train the decision tree. During run-time, the decision tree classifies a new configuration \( V_u \) to some value \( R_{\text{max}}^u \), which allows admission control to use Eq. (1) to obtain the performance and decide whether the configuration is accepted.

6. Experimental Results

To validate our admission control mechanism, we 1) test the accuracy of the \( R_{\text{max}} \) values predicted using the decision tree, which ensures that the admission control mechanism correctly evaluates a configuration, and 2) measure the performance of the decision tree to assess its suitability for responsive systems. We perform a number of tests using a publish/subscribe implementation of the real-time middleware described above, programmed in Java\(^1\). The tasks are publishers (information producers) producing a rate of events (messages on a topic) that are sent over the network to subscribers. Each publisher uses a different topic. Other nodes run tasks that are subscribers (information consumers) on those same topics, which we use to measure the rates without altering the execution of the publishers on the first node. Each machine is an IBM® HS20 blade with 2 Intel Xeon\(^2\) 3.20GHz 64-bit processors with Hyperthreading\(^{\text{TM}}\) enabled and 2 Gbyte of RAM. The machines are connected directly using a Gigabit Ethernet switch. The message size is 128 bytes, so that even at the highest sending rate the network bandwidth is not an issue. Confidence intervals are 95%. The observed range for \( R_{\text{max}} \) is 60,000–100,000 msg/sec, for an attempted total sending rate of 100,000 msg/sec.

**Parameter Estimation Error.** In a real system, we expect to find relatively complex configurations. However, the number of possible configurations explodes as they get more complex. Thus, we expect a system administrator to know in which regime the system is supposed to operate, providing an indication of which configurations are the most common around which the system works.

![Figure 3. Comparing the error estimating \( R_{\text{max}} \) between the decision tree and the average.](image)

To validate our approach, we use a fraction \( C_{\text{train}} \) of the configurations in \( C \) to train the decision tree, and use a disjoint set \( C_{\text{test}} \) to measure the accuracy of the estimated \( R_{\text{max}} \) (i.e., the output class of the decision tree) compared with the actual \( R_{\text{max}} \) obtained during the calibration. Most of the time, the decision tree does

1. Java is a trademark of Sun Microsystems, Inc. in the United States, other countries, or both.
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Figure 4. Performance of the C4.5 decision tree algorithm estimating $R^\text{max}$ from configurations.

not provide a unique possible classification for a given configuration. A number of possible classes $R^\text{max}_i$ are provided with a probability $p_i$ associated with them. We pick the class with the highest probability as result. In an attempt to obtain better estimations, we also try the weighted average of the possible classifications, $\sum_i p_i R^\text{max}_i$. We compare the error estimating $R^\text{max}$ between the decision tree and a naive predictor that takes the average of the parameter values in $C_{\text{train}}$.

Each class spans a range of 1,000 msg/sec. First we use 10 publishers and run $|C|=1,000$ different configurations. We generate 10 starting configurations where 1 publisher gets all the resources, 2 publishers get all the resources equally, etc., and we generate 100 random variations from each of them. This setup is close to a system administrator providing a number of commonly used configurations, and extending the training set by adding similar configurations around the points of operation. The results are shown in Fig. 3, where we train the algorithms with 100 to 900 configurations in $C_{\text{train}}$ and test them with another 100 chosen at random in $C_{\text{test}}$. The naive approach produces much higher error, showing the variation in the values that the model parameter $R^\text{max}$ can take (otherwise the average would be an excellent predictor). The decision tree produces an error in the estimation very close to the quantization error. Also in this case, testing for many configurations does not seem to pay off. The average of the tree decision output does not improve the quality of the result (see plot Decision Tree (avg) in Fig. 3). Then we try a configuration with 20 publishers, which doubles the number of degrees of freedom in the test. We run the experiment for 2,000 configurations generated like the ones in the previous experiment, but starting with 20 configurations. $C_{\text{train}}$ contains from 100 to 1,900 configurations, and we still test with another 100 configurations. With the larger configuration, the error estimating $R^\text{max}$ is close to 700 msg/sec, which is worse but still close to the quantization error. Still, results are very similar to those in Fig. 3. Running 200–300 configurations in $C_{\text{train}}$ seem to provide all the benefit that the algorithm can bring; so again, even for fairly large configurations the calibration phase should not take long. All the following experiments use this last, more complex configuration (20 tasks).

Class granularity. We perform the same tests but using a tree with classes spanning a shorter range: we divide the output into 500, 200, 100 and 10 msg/sec classes, which make for 80, 200, 400 and 4,000 output classes, respectively. In Fig. 4(a) we show the average improvement of the $R^\text{max}$ estimation measured in msg/sec (more is better). If we look at the results when training the tree with the first few hundred configurations, we see that only the classes of 200 msg/sec provide a significant advantage over the 1,000 msg/sec classification. We observe though, that when we increase the configurations in $C_{\text{train}}$ above 1,000, all classifications improve their accuracy over the 1,000 msg/sec classes. Reducing the class size to 500, 100 or 10 improve the estimation by close to 200 msg/sec; the improvement is close to 250 msg/sec when using 200 msg/sec classes. Using 10 msg/sec classes unnecessarily increase the complexity of the tree with 4,000 possible outputs, and yields performance similar to the other classifications.

Computational cost. Decision tree algorithms in general, and C4.5 in particular, are computationally efficient algorithms. In our approach, this computation is part of the calibration phase of the system, an off-line process where responsiveness is not an issue. However, it may be useful to recalculate the decision tree during the on-line operation of the system, as more and more configurations are being used and the resulting parameter values can be gathered. This assumes that the system changes its configuration on a fairly frequent basis. Fig. 4(b) shows the time needed to compute the decision tree using different sizes of $C_{\text{train}}$ and
different class sizes, with C4.5’s C implementation, which can be called from the middleware using the Java Native Interface (JNI), on a Red Hat®/Linux® system with an Intel Pentium IV CPU at 2.66 GHz. We observe that for the initial classification in 40 classes of 1,000 msg/sec, the time needed to compute the decision tree is under 100 ms, using up to 300 configurations. As we have previously seen, more configurations do not bring much benefit to the precision of the parameter estimation. We can then realistically recompute the decision tree on-line with new configurations gathered during operation. Increasing the number of classes increases the precision of the estimation, but only if we train the tree with enough classes (Fig. 4(a)). This brings the computation time into the range of a few hundred millisecond. Fig. 4(c) shows the time taken to query the decision tree with a configuration of 20 tasks. The time remains more or less constant, independently of the number of configurations used to train the decision tree. The number of classes has a sub-linear influence on the query time. Excepting the smallest class size (10 msg/sec), all classifications need around 100–150 μs per query. These figures show that admission control should not affect significantly the operation of the entire system.

7. Conclusion

The upcoming distributed responsive and real-time systems require the allocation of resources to tasks in a complex environment where it is not easy to predict which performance will result from a given resource partition. On a CPU allocation middleware, we propose an admission control mechanism based on a calibration phase prior to deployment of the system in production, where machine-learning algorithms build a knowledge system that is used at run-time to estimate the parameters of the system model for any configuration, enabling the prediction of the system’s performance. In particular, we have shown experiments with decision trees generated with the C4.5 algorithm. We conclude that, by training the system with a few hundred representative configurations, the model parameters can be estimated to a high degree of accuracy (error<1%).

References


3. Linux is a registered trademark of Linus Torvalds in the United States, other countries, or both.


