Actor-based Runtime Model of Adaptable Feedback Control Loops (Position Paper)

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ABSTRACT

Engineering self-adaptive systems is a particularly challenging problem. On the one hand, it is hard to develop the right control model that drives the adaptation; on the other hand, the implementation and integration of this control model into the target system is a difficult and an error-prone activity. Models@runtime is a promising approach to managing adaptations at runtime, as they provide higher levels of abstractions of both the running system and its environment. However, recent work mainly focuses on runtime models that are causally connected to running systems and less attention is paid to how models can be used to develop and manage the control logic that drives runtime adaptations.

In this paper we propose an alternative form of models@runtime as a reactive data-driven model centered around feedback control loops. Both the target system and the adaptation logic are represented as networks of message passing actors. Each of these actors represents a particular abstraction over the running system (sensors, effectors) and its control (analysis, decision). Moreover, the actors are also viewed as target systems themselves. This makes the feedback loops adaptable at runtime as well and permits us to build complex solutions with hierarchical layers of control loops. We discuss how this representation fits some of the requirements of models@runtime and helps to prototype a feedback control system on a concrete example extracted from ongoing validation case studies.

1. INTRODUCTION

Being able to dynamically adjust software structure and behavior at runtime is increasingly becoming an important aspect of contemporary software systems. Such self-adaptive systems are characterized by the ability to autonomously reconfigure themselves based on their perception of their state and their environment, usually following some higher level user specified goals [3]. While these properties can make systems more robust and prone to failures, their construction is a challenging engineering problem [2].

A key aspect of engineering such systems is the Feedback Control Loop (FCL) as it provides the generic mechanism for self-adaptation [2]. A typical FCL consists of distinct phases of monitoring and analyzing both the running system and its environment, followed by planning an subsequent execution of changes that alter the system state. Self-adaptive systems based on feedback control have been proposed, but their engineering is far from mature [17, 2]. Possible reasons are the difficulty of creating a control model [15, 10], and designing and implementing of the control mechanisms. The integration of control mechanisms into a target application is also a difficult and error-prone activity.

One of the promising approaches supporting development of such systems involves the use of models@runtime. This research aims at providing abstractions that facilitate automated management of runtime behavior [6]. The runtime models provide higher-level abstractions of both the running system and its environment and are also manipulated at runtime. However, the focus has been mostly on effecting changes through modifications to runtime models that are causally connected to running systems. Less attention has been paid to how runtime abstractions can be used to develop and manage the control logic that drives runtime adaptations.

In this paper we propose an alternative form of models@runtime in which the model is a reactive data-driven model of feedback control loops. The model is based on a runtime architecture that separates the target of adaptation (e.g., a software system that can be observed and/or modified) from the control layer. Actors [11] are used to represent elements in the target system (e.g., sensors and effectors) and the control layer (e.g., analysis and decision-making mechanisms). It consists of configurations (networks) of actors representing the target system and control layers. Different forms of control can be obtained by linking actors in different ways. The actor representation makes the FCL explicit through models at both design and run times. This explicit representation provides abstractions that should make it easier to reason about the system and the correctness of interactions between different FCL elements. This is one of the main benefits of using a model to describe adaptation processes [21].

Moreover, by introducing reflective capabilities into this model we are able to observe and modify the actors and their configurations uniformly, in the same way as the running system. This makes the FCLs adaptable at runtime as well as the target system and thus enables one to build complex solutions in cases when there is a need for control coordination or hierarchical layers of control loops.

The rest of the paper is organized as follows. Section 2 gives background on the actor model. Section 3 presents an overview of our approach to constructing feedback control loops using the actor model. It is followed by an illustration of this approach on a concrete example (Section 4). Section 5 positions our work with regards to the related work. Finally we discuss how this representation fits models@runtime and its requirements together with an
outline for further work in Section 6.

2. BACKGROUND

The actor model provides a concurrent programming paradigm that is based on actor objects [11]. Actors are essentially independent concurrent processes that encapsulate their state and behavior and communicate exclusively by exchanging messages. Messages are sent asynchronously and each actor maintains a queue of received messages (a mailbox) and processes them one-by-one.

Actors act independently of one another. In its core it is a simple entity that can only receive, process and send messages. The actor states are stored within the actors themselves (states are not globally accessible). An actor can only manipulate its own state and never directly accesses the state of other actors, however, it can send a message that in turn may change the state of its recipient.

The model is inherently reactive. Each actor waits for messages in order to do its work. The task an actor performs upon message arrival is completely left for the user to specify. Different behaviors can be associated to different messages and it can also include spawning new actors.

Actors are independent of threads. Different implementations can map threads to actors in different ways. It is possible to have more than one thread to execute one actor and similarly to have only one thread executing multiple actors. However, the actual message processing is always executed by one thread at the time and the messages must be immutable in order for the model to work correctly.

Following are the most relevant features of actors in our context:

1. Thread safe. The actor model provides a notable advantage over the standard shared memory concurrency. Accessing an actor’s mailbox is by design race condition free. A lock free code is also simpler to write and maintain. It is important to note that actors are not dead lock free as it is possible to have two actors to wait for messages from each other.
2. Dynamic. Actors can create new actors.
3. Scalable. The actor model has a lower context-switching overhead over the standard shared-memory threads with locks, as well as a reduced memory footprint [8].
4. Distributed. The message-based concurrency with encapsulated state can be seen as a higher-level model for threads with the potential to generalize to distributed computation [8].
5. Available. Today, there exists several high-performing actor libraries for mainstream programming languages1.

3. ACTOR REPRESENTATION OF FEEDBACK CONTROL LOOPS

In this section we present an overview of our approach that will be followed by a concrete example in the next section.

3.1 Overview

Our approach is centered around the concept of Feedback Control Loop (FCL) and we use the actor model to represent it. Each actor in the model has one of the following roles that correspond to some aspects of a FCL.

Sensor. Collects raw information about the state of the system and its environment from observable entities like various operating system resources, service calls, user preferences, etc. A sensor can be either active or passive. An active sensor is used when the data update is based on some external event, like a file change or a timer.

Effectors. Carries out changes on the target system. It can be seen as a parameterized operation that is executed in order to alter the state of the system.

Controller. Represents the decision making process. Essentially, it is responsible for choosing the appropriate action(s) among the set of all permissible actions based on the state of the running system that has been inferred by sensors and filters. There are many different types of controllers that can be employed. A good overview of some of the different control strategies can be found in [15].

The first two roles are related to the relevant abstractions over the running system and its surrounding environment (touch-points), the two others represent the adaptation logic.

The actual adaptive behavior is realized by connecting together corresponding actors (cf. Section 4.2) to form an acyclic directed network. It can be partitioned into layers: a system layer consisting of sensors and effectors that provide all the inputs and outputs necessary for the adjacent control layer to reason about and to adjust the execution context of the running system represented by the system layer.

The different concerns of FCLs are separated and realized by one or more collaborating actors. The (1) data collection is realized by sensors, (2) analyzing is done by filters, (3) decision is taken by controllers, and (4) finally effectors are responsible for final changes execution. Since actors are independent entities, some of them could be reused in different systems. The primary motivation of this separation of concerns is to allow for building the loops from composed rather than programmed entities, as well as to foster reuse of both the independent actors and partial or complete loops.

3.2 Model Reflection

The reflection support in the model is also realized using actors. We provide actors that represent the introspection and intercession capabilities of the actor runtime system. This includes effectors controlling the sensors life-cycle allowing to start new actors or stop the already running ones, changing message recipients, as well as active sensors sending notifications every time an actor changes its state (cf. Section 4.5).

Conceptually we can view any actor as a running system itself. We are therefore able to observe and modify its state using touch points similarly as with the running system. We use a special notion of provided sensors and provided effectors to implement this actor reflection. Technically, they are just ordinary sensors and effectors actors that use special messages to indirectly manipulate the state of other actors and together form another system layer of the control system (cf. Figure 3). They are logically bound to the context of other actors, parents, that are responsible for managing their life cycle.

Provided sensors are inherently active. They are notified by their parents every time their state has changed so the information can be further pushed into the control layer. This is important as it allows for assembling the control layer in a uniform way regardless of what system layer is being use.

1http://goo.gl/WcXCE
This actor reflection is a crucial feature that permits to build complex solutions in cases when there is a need for control loops coordination or for organizing them into hierarchical layers. The coordination can be done by explicitly synchronizing controllers via provided sensors (cf. Section 4.4). The hierarchical layering uses reflection to enable loops in one layer to be supervised by another loop in the adjacent higher-level layer (cf. Section 4.3).

Since both provided sensors and effectors are actors themselves, they can also be observed and modified using their provided sensors and effectors. However, in our studies and validations, we do not find more than two levels to be very useful.

3.3 Connection to the Running System

The connection between the touch point actors, i.e., sensors/effectors, and the running system is realized using uni-directional causal links that use the introspective (sensing) and intercessory (effecting) capabilities of the provided system interfaces. While these abstractions might seem to be too low-level [19], they simplify the actual implementation of the execution code and system reasoning. Filters are employed to infer higher-level information from the sensor data. This way we can provide flexible abstractions based on the need of a particular controller for specific adaptation scenarios.

From an implementation perspective, the actual actor behavior, i.e. the executable code, is specified through a delegate that decouples the actor related code from the user code. In practice this usually means creating a new class that complies to a certain interface depending on the target actor role. Since there is no need for any actor specific code in the delegate, it can be unit tested in isolation. It is also important to note that even though the delegate code might be used in highly concurrent environments, it does not have to be thread safe as it is only called by one thread at a time. This considerably simplifies the coding efforts required from the final developers.

3.4 Execution

The loop execution is driven by messages that are passed among actors. By design the actors are passive, they sit waiting for messages to arrive. We use active sensors as a triggering mechanism for loop execution.

During its initialization, an active sensor registers itself to some external events. When an event occurs, it extracts the relevant information and dispatches it to the corresponding actors (i.e., filters or controllers). This in turn initiates other messages before the actual control code is executed. An example of the messages passing is given in the sequence diagram schema on Figure 2.

Consequently the reactive nature of the model results in less need for synchronization between the running system and the abstraction layer, since the introspection and intercession happen during the loop execution. However in some cases, implicit synchronization is inevitable. It is thus supported in the model (cf. Section 4.5).

4. ILLUSTRATIVE CASE STUDY

4.1 Motivating Adaptation Scenario

We consider a High-Throughput Computing environment where users can log in an interface allowing them to execute scientific workflows using the Condor infrastructure [18]. Condor is a well-established distributed batch computing system that has been extensively used in both academia and industry.

In Condor, a workflow is usually executed by running a DAGMan that takes control of the jobs dependencies and submits the ready to be executed jobs to the Condor scheduling agent (schedd). The schedd is responsible for managing the queue of the user submitted jobs and mapping them onto a set of resources where the actual execution is performed. By default schedd accepts all valid submissions requests regardless of the current state of the system. Since especially scientific workflows tend to be large, it can become a bottleneck and potentially overloaded as the work it has to perform is proportional to the number of jobs that it controls. This in turn might result into an overall throughput degradation. There are a number of configuration options for both the scheduler and the workflow manager that allows for fine tuning the system behavior. However, all of them are statically defined, without considering the current state of the system.

Our objective is therefore to graft a control on top of the Condor job submission, which would ensure high throughput, while preventing the schedd overload. One way is to base the adaptation on a model that takes the number of jobs in the queue (N), together with a service rate (µ) and a number of executing DAGMans (m) to compute a delay (d) that is used to throttle workflow submission rates. This will help to keep the number of idle jobs within a certain range for which the system performs well. More information about the control model together with some experimental results are given in [13].

4.2 Modeling Feedback Control Loops

Figure 1 illustrates the actors and the message flow. It includes the causal links to the running system of one feedback control loop that implements the illustrative example using the control elaborated above.

The control loop is triggered by the trigger sensor that periodically (every p seconds) notifies the connected controller by sending an arbitrary signal s. Before the controller can compute the delay d, it needs to gather required data from the connected sensors:

1. The $queueStats provides the information about the number of jobs in the queue N using the $condor_q command.
2. the current service rate µ comes from the condor history file and is extracted by the serviceRate, and
3. the number of running DAGMan processes \( m \) is obtained from
the \texttt{dagmanCounter} by executing the system \texttt{ps} command.

In order to increase the stability of data and avoid oscillation, the
\( N \) and the \( \mu \) inputs are filtered using moving average filters before
they enter the \texttt{submissionRateCtrl}. We do not need to filter
the number of DAGMans \( m \) since workflow executions are usually
long running processes. Once all the inputs have been gathered
the controller, using the formulas elaborated in [13], computes the
delay \( d \) that will be imposed system wide to all running DAGMans
by the \texttt{dagmansDelayer}. In this example it is simply written
into a file from which it is read by DAGMan the next time it tries
to submit a job.

While some of the actors identified are use-case specific such as
the \texttt{submissionRateCtrl}, others could be reused in some dif-
ferent Condor related control systems. One of our ongoing long
term goal is to initiate a library of reusable FCL elements and pat-
tterns, both at the model and implementation levels.

![Figure 2](image.png)

Figure 2: Excerpt of the message sequence for the control ex-
ample from Figure 1. For simplification, we omit the filters. The
delegate() call represents the execution of the actual user code.

The Figure 2 shows a translation of the actor schema into a se-
quence diagram representing the actual message passing chain. The
messages are all sent asynchronously. This is thus the only in-
stance of the possible execution sequence. The \texttt{submission-
RateCtrl} can receive responses from the connected sensors in
different order.

### 4.3 Hierarchy of Feedback Control Loops

The primary concern in the above overload control was to protect
an unbounded usage of the schedd submission service. The prob-
lem is that the protection of one resource might in consequence
jeopardize another and thus transform one problem into a differ-
ent and potentially much worse one (as it might be unexpected).
Therefore, in general, this process has to be done recursively for
any potential unbounded resource usage, regardless whether it is a
system resource or a service usage.

In our solution this problem is related to the \texttt{queueStats} sen-

tor. Internally it executes \texttt{condor_q} command that makes the
schedd walk its entire job queue, which becomes an expensive op-
eration when the number of jobs is large (this behavior is internal
to Condor). One way to fix this is by adding a new control layer on
top of the existing one that will throttle the execution frequency of
the \texttt{queueStats} sensor.

The new control schema is depicted in Figure 3. First we extend
\texttt{CondorQueueStats} to provide information about how long it
took the last time to execute \texttt{condor_q} using a provided sensor
\texttt{execTime}. Next we add a provided effector \texttt{setPeriod} to the
\texttt{PeriodicSignal} sensor that allows for adjusting the triggering
period \( p \). Finally, in the new meta-control layer we create a con-
troller that adjusts the execution frequency proportionally to the
\texttt{execTime} time.

### 4.4 Coordination of Feedback Control Loops

Actor reflection is also used for coordinating multiple interacting
loops using an explicit synchronization. For example when one
controller reaches a certain state, it can use a provided sensor to
signal it to another controller that then performs its own work.

The following figure 4 shows how a possible coordination sce-
nario can be implemented from our overload control use case.

![Figure 4](image.png)

Figure 4: Example of feedback control loops coordination. The \( \emptyset \)
denotes no notification.

The \texttt{submissionRateCtrl} might observe that the service
rate measures are somehow corrupted and explicitly requests a re-
pair loop to be executed. It notifies the provided sensor \( S_1 \) which
will indirectly trigger the \texttt{resetHistory}, which will try to fix
the situation.

Like with any other synchronization, some of the technical de-
tails have to be handled, for instance: What should happen if the
repair loop takes too long to execute and the \texttt{submissionRate-
Ctrl} has executed in the meantime, still finding the value to be
corrupted? What will happen if, by the repair action, the service-
Rate sensor crashes? Since often the solutions to these issues are
case specific, the model itself does currently not provide any spe-
cial support. They have to be resolved manually. However, we are
currently investigating how one can express these concerns more
explicitly.

### 4.5 Implicit Synchronization

In the presented control system we use one actor that represent
all the running workflow managers. When a new delay is computed
and pushed it affects all running DAGMans. Since we use only one global delay, this is a correct abstraction. However, we can imagine to have a more complex control that will compute different delays for different DAGMans, therefore needing to represent each running DAGMan explicitly as an actor.

Figure 5: Example of an implicit synchronization of running DAGMan processes

Figure 5 depicts a control loop that can be used for an implicit synchronization between a running system and its runtime abstraction. The newProc sends out a notification every time a new process has been started. In the filter DAGManProcFilter, processes that are not DAGMans are filtered out and the others are pushed to a controller that uses the actorDeployer to spawn a new actor of an appropriate role, type, and configuration. Of course a similar loop needs to be built for stopping actors whose underlying DAGMan has finished execution.

5. RELATED WORK

IBM proposed a form of standardized approach for the feedback loops with the Monitor-Analyze-Plan-Execute-Knowledge (MAPE-K) architecture [12]. Several systems and frameworks have been developed according to this principle. For example, the Rainbow architecture [7] consists of two-layer framework with an external fixed control loop. While the loop is made explicit, the framework is not explicitly reflective and therefore it does not support hierarchical FCLs.

The Genie domain specific modeling tool [1] uses architectural models to support generation and execution of adaptive systems for component-based middleware technologies. The adaptive logic is specified as state machines. Similarly, the DiVA project [16] also uses architectural models for dynamic adaptation, together with modeling some other aspects such as variability, system context and adaptation logic. It employs MDE techniques to produce re-configuration scripts that performs the actual adaptation. In our approach we rather focus on a more general and explicit feedback control loop modeling framework.

In [20] Vogel et al. promotes the use of runtime executable megamodels. They present a modeling language for adaptation logic model together with a runtime interpreter that executes the megamodels. Similarly to our approach they can also represent loop coordination and hierarchically organize them into layers. However, while they provide illustration of coordination and layering of different adaptation strategies, there is only a high-level overview of how the actual adaptations look like. On the contrary our model provides a concrete runtime representation of a feedback control loop. Their solution relies on an implicit synchronization between the runtime megamodels and running system. The metamodel is based on EMF that has some limitations for the use at runtime such as higher memory footprint, lack of thread safe access to the model [5].

In [9] Hebig et al. provide a UML profile to explicitly architect several coordinated control loops with component diagrams. Nevertheless the loop elements are not adaptive and it does not explicitly show the adaptation flow. There is also no connection to a runtime support.

As for control, in [14], Litou et al. describe a hierarchical framework that accommodates scopes and timescales of control actions, and different control algorithms. Their architecture considers three main types of controllers reflecting the three different stages that they focus on: tuning, load balancing, and provisioning. While being similar, our architecture is more general, but provides less fine-tuned building blocks to control the behavioral models used inside controllers.

There is also a large body of work that concerns designing feedback control for embedded computing, for example Ptolemy II [4]. We have a similar actor-based approach whose execution semantics is comparable with Ptolemy’s component interactions model of computation. We need to further investigate the main assumptions and differences between these systems so to determine the possibility of reuse, i.e. being able to generate Ptolemy II model from our actor model.

Similar approach could be realized using more traditional reflective component model. However, in such a case next to the component styles that would correspond to our actor roles, one would also have to define an execution semantics (the component states, state’s transitions and model of computation).

6. DISCUSSION

In this section we discuss how our proposal fits in the landscape of models@runtime and the overall goal of FCL engineering for self-adaptive software systems. We also give an overview of ongoing and further work.

6.1 Towards Models@runtime

Even though it is not a typical approach to models@runtime we believe that the proposed solution is a possible and pragmatic approach towards realizing feedback control systems. The actor model provides a higher-level abstraction of a running system that is manipulated at runtime to dynamically adapt the underlying system. The concern is whether the proposed abstraction is expressive and intuitive enough to support activities of runtime software adaptations based on feedback control.

To validate this point we need to conduct more experiments and design self-adaptive systems across different domains. However, the preliminary results already show the following major benefits that considerably ease development: (1.) it provides an explicit and concrete representation of FCLs described in the actor model, (2.) the adaptation in form of FCLs is specified at a higher-level of abstraction by a network of cooperating actors, (3.) each actor is an independent unit that can be used in different adaptation contexts, (4.) the reflection capabilities of the model allows the actors to be adapted at runtime in a uniform way, and (5.) multiple coordinated and hierarchically layered FCLs can be expressed in a uniform way.

Moreover, we find actors actor to be a particularly suitable model, as it directly supports some of the main non-functional requirements of models@runtime [5]. Thread safety, offered by design, results in potentially more secure and easier to implement execution code. Since the actors encapsulate their state and communicate only by sending asynchronous messages, they are suitable for distributed environments. The actor model is designed to be scalable and lightweight. For instance in the Akka® 2.0 the memory overhead is about 400 bytes per instance (2.7 million actors per GB of heap) with a possible throughput of 50 million messages per second.

2http://akka.io
sec on a single machine\textsuperscript{3}. This combined with many actor model implementations available today enable one to deploy the proposed solution on top of a significant range of systems.

6.2 Towards Rapid Prototyping

The model elaborated in this paper is now being integrated into our Model-Driven Engineering toolkit ACTRESS\textsuperscript{4}. Its aim is to use modeling techniques to provide engineers and researchers with a tool to build models that help them to rapidly prototype and experiment with self-adaptive systems, putting quickly their ideas into practice. The approach needs to be built on a flexible abstraction and must cope with the diversity and heterogeneity of current platforms, while being efficient in performance.

The system is currently being implemented in Scala\textsuperscript{5} using the aforementioned Akka framework. From the initial results the actor model seems a suitable abstraction that performs well, but empirical evidence and wider usage are still missing. Besides it must be noted that the model itself is technologically agnostic. Other languages and frameworks can be used to target different classes of systems.

6.3 Ongoing and Future Work

The proposed approach is currently under validation inside the ANR-funded SALTY project\textsuperscript{6}. Ongoing work on this project use-cases (e.g., regulation of a geo-tracking system following fleets of thousands of trucks) will complement validation by providing more thorough experience of applying the proposed approach. We are currently in the process of making the implementation usable with a comprehensive set of unit tests over the modelling support. However the complete MDE tool chain is still at the prototype level.

Further, we need to work on the composition of actors and how to express some of the concerns related to loop coordination and failure propagation.

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7. REFERENCES


