A Novel OFDM PAPR Reduction Scheme Using Selected Mapping Without Explicit Side Information

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Abstract—Selected mapping (SLM) is a technique used to reduce the peak-to-average power ratio (PAPR) in orthogonal frequency-division multiplexing (OFDM) systems. SLM requires the transmission of several side information bits for each data block, which results in some data rate loss. These bits must be channel-encoded because they are particularly critical to the error performance of the system. This increases the system complexity and transmission delay, and decreases the data rate even further. In this paper, we propose a novel SLM method for which no side information needs to be sent. By considering the example of several OFDM systems using either QPSK or 16-QAM modulation, we show that the proposed method performs very well in terms of bit error rate at the receiver output provided that the number of subcarriers is large enough.

I. INTRODUCTION

High peak-to-average power ratio (PAPR) is a well-known drawback of orthogonal frequency-division multiplexing (OFDM) systems. Among all the techniques that have been proposed to reduce the PAPR, selected mapping (SLM) is one of the most promising ones because it is simple to implement, introduces no distortion in the transmitted signal, and can achieve significant PAPR reduction [1]. The idea in SLM consists of converting the original data block into several independent signals, and then transmitting the signal that has the lowest PAPR. The selected signal index, called side information index (SI index), must also be transmitted to allow for the recovery of the data block at the receiver side, which leads to a reduction in data rate. This index is traditionally transmitted as a set of bits (the SI bits). The probability of erroneous SI detection has a significant influence on the error performance of the system since the whole data block is lost every time the receiver does not detect the correct SI index. In practice, a channel code must thus be used to protect the SI bits. This further reduces the data rate, makes the system more complex, and increases the transmission delay.

It may therefore be worth trying to implement the SLM method without having to explicitly send any SI bit. A few techniques for doing so have already been proposed such as, for example, the scrambling method described in [2] or the maximum-likelihood decoding scheme introduced in [3]. More recently, another method was proposed in [4] that combines channel estimation using high-power pilot tones and PAPR reduction via SLM. One of the key ideas in [4] consists of choosing the location of the pilot tone used inside each data sub-block depending on the SI index, and then exploit the power disparity between this pilot tone and the data symbols in the same sub-block in order to recover this index at the receiver side. The technique in [4] assumes the use of several pilot tones in each data block, which is not a suitable option for fading channels that change slowly. The practical limitations of the work in [4] can however be suppressed by leaving out the pilot tones and replacing them with data symbols. By doing so, we remove the restrictions regarding the possible locations of the high-power subcarriers inside the data block, and are therefore left with a problem which is much more general than that addressed in [4]. In this paper, we address this problem by describing a novel SLM method without side information and studying its performance in terms of probability of erroneous SI detection and bit error rate (BER).

II. THE PROPOSED SLM TECHNIQUE WITHOUT SIDE INFORMATION

In this section, a notation in the form $V = (v_q)_Q$ shall be used to denote a vector $V$ composed of $Q$ scalar quantities $v_q$, $q \in \{0, 1, ..., Q - 1\}$.

A. The proposed SLM transmitter

Consider an OFDM system using $N$ orthogonal subcarriers. A data block is a vector $X = (x_n)_N$ composed of $N$ complex symbols $x_n$, each of them representing a modulation symbol transmitted over a subcarrier. In classical SLM, $X$ is multiplied element by element with $U$ vectors $B_u = (b_{u,n})_N$ composed of $N$ complex numbers $b_{u,n}$, $u \in \{0, 1, ..., U - 1\}$, defined so that $|b_{u,n}| = 1$, where $|\cdot|$ denotes the modulus operator. Each resulting vector $X_u = (x_{u,n})_N$, where $x_{u,n} = b_{u,n} \cdot x_n$, produces, after inverse discrete Fourier transform, a corresponding vector $X'_u$ composed of $N$ symbols $x'_{u,q}$, $q \in \{0, 1, ..., N - 1\}$, given by

$$x'_{u,q} = \frac{1}{\sqrt{N}} \sum_{n=0}^{N-1} x_{u,n} \cdot \exp \left( j \frac{2\pi u q}{N} \right).$$  \hspace{1cm} (1)

Among the $U$ vectors $X'_u$, the one leading to the lowest PAPR, denoted as $X'_{v}$, is selected for transmission. To allow for the recovery of the original vector $X$ at the receiver side, one needs to transmit the index $v$ of the selected vector. This index, hereafter called SI index, is generally transmitted as a set of $\log_2(U)$ bits. In this paper, we propose a novel SLM technique that allows for the SI index $v$ to be reliably embedded in the transmitted vector $X'_v$ so that no additional bits need to be sent to the receiver. In this technique, the vectors $B_u = (b_{u,n})_N$, $u \in \{0, 1, ..., U - 1\}$, are such that, for example, the scrambling method described in [2] or the maximum-likelihood decoding scheme introduced in [3]. More recently, another method was proposed in [4] that combines channel estimation using high-power pilot tones and PAPR reduction via SLM. One of the key ideas in [4] consists of choosing the location of the pilot tone used inside each data sub-block depending on the SI index, and then exploit the power disparity between this pilot tone and the data symbols in the same sub-block in order to recover this index at the receiver side. The technique in [4] assumes the use of several pilot tones in each data block, which is not a suitable option for fading channels that change slowly. The practical limitations of the work in [4] can however be suppressed by leaving out the pilot tones and replacing them with data symbols. By doing so, we remove the restrictions regarding the possible locations of the high-power subcarriers inside the data block, and are therefore left with a problem which is much more general than that addressed in [4]. In this paper, we address this problem by describing a novel SLM method without side information and studying its performance in terms of probability of erroneous SI detection and bit error rate (BER).
each $B_u$, the moduli of $K$ elements $b_{u,n}$ are set to a constant $C > 1$, called *extension factor*, whereas the moduli of the other $(N - K)$ elements $b_{u,n}$ remain equal to the unit. Note that the phases of these elements can be set to any random values as in classical SLM. For a given vector $B_u$, the locations of the elements $b_{u,n}$ for which $|b_{u,n}| = C$ form a set $S_u$ composed of $K$ integers. For instance, if a vector $B_u$ is associated with the set $S_u = \{0,7,25,37\}$, it means that only the complex elements $b_{u,n}$ in positions $n = 0,7,25,$ and $37$ in $B_u$ have a modulus equal to $C$. The set $S_u$ is actually representative of the index $u$. To allow for SI recovery in the receiver, there must be a one-to-one correspondence between an index $u$ and a set $S_u$. In other words, two distinct vectors $B_u$ must NOT be associated with identical sets. The number of distinct sets $S_u$ that can be generated is given by the binomial coefficient \( \binom{N}{K} \). Since the SI index can take $K$ values, we only need to use $U$ sets $S_u$ among the \( \binom{N}{K} \) available sets.

Once all vectors $B_u$ have been generated, our SLM method works exactly like the classical one, i.e. the data block $X$ is multiplied by element with vector $B_u$ so as to produce $U$ vectors $X_u = (x_{u,n})_N$, with $x_{u,n} = b_{u,n} \cdot x_n$, as well as $U$ corresponding vectors $X_u'$. In a given vector $X_u$, the symbols $x_{u,n}$ with $n \in S_u$ have an average energy $C^2$ times greater than that of the other symbols, and are hereafter referred to as *extended symbols*. This disparity in average energies between extended and non-extended symbols is what allows the receiver to recover the SI index after transmission. Finally, the vector $X_u'$ with the lowest PAPR is transmitted. We recall that, throughout this paper, this particular vector is denoted as $X_u'$ and is associated with the vectors $X_v = (x_{v,n})_N$ and $B_v = (b_{v,n})_N$. The index $v$ represents the SI index to be transmitted.

In the proposed SLM technique, the average energy per transmitted symbol is increased when the data block $X$ is multiplied by the vectors $B_u$ because the fact that $|x_{u,n}| = |b_{u,n} \cdot x_n|$, with $|b_{u,n}| = 1$ or $C$, implies that $E[|x_{u,n}|^2] > E[|x_n|^2]$, where $E[\cdot]$ designates the expectation operator. Note that any energy increase in $X_u$ translates into an identical energy increase in $X_u'$ (Parseval’s theorem). This energy increase $G$, expressed in decibels (dB), is given by

$$G = 10 \cdot \log_{10} \left( 1 + \delta \cdot (C^2 - 1) \right),$$

where $\delta = \frac{C^2}{1}$. In practice, the energy increase is compensated for by decreasing the average energy per complex symbol $x_n$, i.e. reducing the minimum Euclidean distance between signal points in the constellation from which the symbols $x_n$ are drawn. Remarkably, this does not necessarily result in some significant error performance degradation at the receiver output, as will be seen later.

### B. The proposed SLM receiver

In this paper, we assume transmission over a quasi-static frequency-selective Rayleigh fading channel with $Z$ equal-power taps. For each transmitted symbol $x_{v,q}$, $q \in \{0,1,\ldots,N-1\}$, the corresponding received sample $y_q$ is thus given by

$$y_q = \sum_{z=0}^{Z-1} h_z' \cdot x_{v,q-z} + n_q',$$  

where $h'_z$ is a complex zero-mean Gaussian sample representing the fading experienced by the $z$th tap. In (3), $n_q'$ denotes a complex Gaussian noise sample with zero mean and variance $\sigma^2 = N_0$, where $N_0$ denotes the one-sided power spectral density of the additive white Gaussian noise (AWGN). We also assume that the $Z$ fading samples $h_z'$ are independent and perfectly known at the receiver side, i.e. perfect channel state information (CSI) is considered. Under these assumptions, we can show that, after discrete Fourier transform, the received sample $y_n$ corresponding to the $n$th subcarrier, $n \in \{0,1,\ldots,N-1\}$, is given by

$$y_n = h_n \cdot x_{v,n} + n_n,$$  

where $n_n$ is a complex zero-mean Gaussian noise sample with variance $\sigma^2 = N_0$ and

$$h_n = \sum_{z=0}^{Z-1} h_z' \cdot \exp\left(-j \frac{2\pi n z}{N}\right)$$

is a complex Gaussian noise sample with zero mean and unit variance. The role of the SI detection block is to recover the index $v$ by processing both vectors $Y = (y_n)_N$ and $H = (h_n)_N$, the latter being computed using (5). This block has the knowledge of the $U$ possible vectors $B_u$ and the one-to-one correspondence between an index $u$ and a set $S_u$.

In this sub-section, we propose an algorithm that exploits the disparity in average energy between the $K$ extended symbols and the $(N - K)$ non-extended symbols in vector $X_v = (x_{v,n})_N$ so as to recover the SI index $v$. Assume a particular location $n \in \{0,1,\ldots,N-1\}$ in the received vector and its corresponding fading sample $h_n$. If the SI index was $u \in \{0,1,\ldots,U-1\}$, the average received energy at location $n$ in the absence of additive noise would be given by

$$E[|h_n|^2|x_{u,n}|^2] = |h_n|^2 |b_{u,n}|^2 \gamma,$$

where $\gamma$ denotes the average energy per complex symbol $x_n$. At the same time, we can show that the average energy of the received sample $y_n$ is given by

$$E[|y_n|^2] = |h_n|^2 |b_{v,n}|^2 \gamma + \sigma^2.$$

At this stage, we can introduce a metric $\alpha_{u,n}$ defined as

$$\alpha_{u,n} = E[|y_n|^2] - \sigma^2 - |h_n|^2 |b_{u,n}|^2 \gamma.$$  

By combining (7) and (8), we can show that

$$\alpha_{u,n} = |h_n|^2 \gamma \cdot |b_{v,n}|^2 - |b_{u,n}|^2 |.$$  

The minimal value of the metric $\alpha_{u,n}$ is equal to 0 and obtained for $|b_{v,n}| = |b_{u,n}|$. Extending this reasoning to the whole received vector rather than focusing on a particular
location \( n \) leads us to introduce another metric \( \beta_u \) defined as
\[
\beta_u = \sum_{n=0}^{N-1} \alpha_{u,n}.
\]
(10)

The minimal value of this metric \( \beta_u \) is also equal to 0 and obtained when the condition \( |b_{u,n}| = |b_{u,n}| \) is satisfied for all values of \( n \), i.e. \( B_v = B_u \) or equivalently \( v = u \). This shows that the SI index can be recovered by determining, using (8) and (10), the value of \( u \) that minimizes the metric \( \beta_u \). It is important to mention that, in the computation of (8), the receiver uses the sample \( y_n \) to approximate the term \( E[|y_n|^2] \) as follows:
\[
E[|y_n|^2] \approx |y_n|^2.
\]
(11)

Evaluating an average by only considering a single sample may obviously lead to a gross estimate of this average. In case we use too many unreliable estimates of the terms \( E[|y_n|^2] \), \( n \in \{0, 1, ..., N - 1\} \), in the computation of (8), the minimal value of the metric \( \beta_u \) may not be obtained for \( u = v \), thus leading to an erroneous detection of the SI index. In order to minimize the probability of occurrence of such an error event, we must ensure that the \( U \) vectors \( B_u \) are as different as possible, i.e. the number of locations \( n \) in two distinct vectors \( B_u \) and \( B_{u'} \) for which \( b_{u,n} \neq b_{u',n} \) is maximised. This should, in most cases, prevent a metric \( \beta_u \), \( u \in \{0, ..., v-1, v+1, ..., U-1\} \), from satisfying the condition \( \beta_u < \beta_u \), even if we have \( \alpha_{u,n} < \alpha_{u,n} \) for several locations \( n \in \{0, 1, ..., N - 1\} \).

III. EXAMPLE

To illustrate the proposed SLM technique, we consider in this section the example of several OFDM systems based on either QPSK or 16-QAM modulation with Gray mapping. Our goal is to achieve PAPR reduction using \( U = 10 \) vectors \( B_u \). All results given in this section are obtained by computer simulations considering the frequency-selective fading channel model previously described with \( Z = 4 \) taps. We also assume the use, at the transmitter output, of a nonlinear solid-state power amplifier (SSPA) simulated using Rapp’s model [5] with a smoothness parameter \( p = 3 \) and an input backoff (IBO) of 7 dB.

A. Construction of the vectors \( B_u \)

In this sub-section, we present a technique to construct the set of \( U = 10 \) vectors \( B_u = (b_{u,n})_N \), \( u \in \{0, 1, ..., U-1\} \). First, for a given vector \( B_u \), the phases of the complex elements \( b_{u,n} \) can be chosen randomly, as already suggested in previous works (see, e.g., [6]). The first step of the procedure used for defining the moduli \( |b_{u,n}| \) consists of dividing the vectors \( B_u \) into subvectors of length \( M \), where \( M \) is the smallest possible integer satisfying the condition \( \binom{N}{k} \geq U \), \( k \) being any integer smaller than \( M \). In our example, \( U = 10 \), and therefore \( M = 5 \) because \( \binom{5}{2} = 10 \) when \( k = 2 \) or 3. If \( U \) was equal to 11 instead of 10, we should have taken \( M = 6 \) because there is no integer \( k \) for which \( \binom{6}{k} \geq 11 \). Each vector \( B_u \) is thus viewed as a succession of \( L = N/M \) subvectors, each of them composed of \( M \) complex elements. Before moving to the next step, we also have to select a particular value of \( k \) among all those satisfying the condition \( \binom{N}{k} \geq U \).

For reasons explained later, we keep the smallest value of the integer \( k \) for further considerations. In our example, we thus find \( M = 5 \) and \( k = 2 \).

Hereafter, we denote by \( b_{m,l} \) the \((m+1)\)th complex element in the \((l+1)\)th subvector, \( m \in \{0, 1, ..., M - 1\}, \ l \in \{0, 1, ..., L - 1\} \). In each subvector, the moduli of \( k \) elements are set to a constant \( C > 1 \), whereas the moduli of the other \((M - k)\) elements remain equal to the unit. The locations in each subvector of the elements \( b_{m,l} \) for which \( |b_{m,l}| = C \) are identical for each subvector, i.e. the moduli \( |b_{m,l}| \) do not depend on the index \( l \in \{0, 1, ..., L - 1\} \). This implies that, once the \( M \) moduli have been defined for the first subvector (corresponding to \( l = 0 \)), this pattern of moduli is simply repeated \((L - 1)\) times in order to obtain the moduli for all locations \( n \in \{M, M + 1, ..., N - 1\} \) in the corresponding vector \( B_u \). The fact that \( \binom{M}{k} \geq U \) implies that there are enough possible permutations of \( k \) extended elements in a \( M \)-element subvector in order to generate the \( U \) distinct sets \( S_u \) that are needed. As an example, Fig. 1 shows the \( U = 10 \) subvectors that can be generated when \( M = 5 \) and \( k = 2 \). Note that Fig. 1 only shows the moduli of the elements \( b_{m,l} \) since their phases can basically take any random values.

The set of \( U \) vectors \( B_u \), obtained using the simple procedure described above is such that, when considering any pair of vectors \( B_u \) and \( B_{u'} \), \( u \neq u' \), the number of locations \( n \) for which \( |b_{u,n}| \neq |b_{u',n}| \) is at least equal to \( \frac{2N}{M} \). In other words, the probability of erroneous SI detection can be lowered by increasing the ratio \( \frac{U}{M} \). This is why it is crucial to minimize, for given values of \( N \) and \( U \), the length \( M \) of the subvectors.

With the construction procedure proposed in this sub-section, the total number of extended symbols in the transmitted vector \( X_u \) is given by \( K = \frac{2N}{M} \). Note that \( \delta = \frac{K}{U} = \frac{K}{N} \). Hence, in order to minimize the energy increase given by (2), we must ensure to use the smallest possible value of \( k \) among all those satisfying the condition \( \binom{M}{k} \geq U \).
algorithm performs better with QPSK than with 16-QAM. This can be explained by the fact that, in QPSK, the energy per symbol $x_n$ is constant, i.e. we always have $|x_n|^2 = \gamma$. This is not the case for 16-QAM where the term $|x_n|^2$ can take three different values: $|x_n|^2 = \frac{\delta}{2}$, $\gamma$, or $\frac{\delta}{4}$, meaning that the term $|x_n|^2$ may significantly differ from its average value $\gamma$. Hence, the corresponding received sample $y_n$, given by (4), has an energy $|y_n|^2$ that tends to deviate more from its average value $E[|y_n|^2]$ in 16-QAM than it does in QPSK. As a consequence, the approximation $E[|y_n|^2] \approx |y_n|^2$ used in the computation of (8) tends to be more accurate in QPSK than in 16-QAM, which makes the occurrence of an erroneous detection event less likely with QPSK than with 16-QAM.

C. Bit error rate performance

It is also important to study the error performance degradation caused by the application of the technique proposed in this paper. Such degradation is potentially due to both the energy increase $G$ and the occasional SI detection error events. In Fig. 3, we have plotted the BER versus $E_b/N_0$ curves obtained for the OFDM systems based on QPSK and 16-QAM, for $N = 65, 125, 255$, and 510 subcarriers. For all simulations, the extension factor value is $C = 1.2$. For comparison purposes, we have also plotted the BER curves obtained with equivalent OFDM systems using classical SLM with perfect, i.e. error-free, side information (SLM-PSI). The latter scenario typically implies the use of a channel code entirely dedicated to the protection of the SI bits during transmission, which is probably more an ideal situation than a practical one.

![Fig. 2. Probability of side information detection error, $P_{de}$, obtained with the proposed SLM technique as a function of the extension factor $C$, for $N = 65, 125, 255$, and 510 subcarriers. The OFDM systems use either QPSK or 16-QAM modulation, with Gray mapping in both cases. The system parameters are: $E_b/N_0 = 10$ dB, quasi-static frequency-selective Rayleigh fading channel with $Z = 4$ equal-power taps and perfect CSI, $M = 5$, $k = 2$, SSPA with $p = 3$ and IBO = 7 dB.](image-url)

B. Probability of SI detection error

Fig. 2 shows the probability of SI detection error, $P_{de}$, as a function of the extension factor $C$, for four different numbers of subcarriers, $N = 65, 125, 255$, and 510, and a signal-to-noise ratio $E_b/N_0 = 10$ dB. We consider two OFDM systems using either QPSK or 16-QAM modulation. For simplicity sake, the numbers of subcarriers are chosen to be multiples of $M = 5$ close to the usual powers of two used in practice. Finally, the algorithm introduced in the previous section is used for recovering the SI index.

From Fig. 2, it is observed that the value of $P_{de}$ depends on the extension factor $C$, the number $N$ of subcarriers, and the modulation scheme that is employed. As $C$ is increased, the performance of the detection algorithm improves simply because a higher value of $C$ allows for a better distinction between extended and non-extended symbols after transmission through the channel, which makes the occurrence of an erroneous detection event less likely. Increasing the number $N$ of subcarriers also results in a lower probability of detection error. We recall that, if we consider two distinct vectors $B_u$ and $B_{u'}$, the minimum number of locations $n$ for which $|b_{u,n}| \neq |b_{u',n}|$ is equal to $\frac{2N}{3M}$. Hence, any increase in the number of subcarriers results in an increase in this number of locations, which then reduces the probability of an erroneous detection event.

The results in Fig. 2 also indicate that our SI detection
of a nonlinear SSPA at the transmitter output. As previously mentioned, throughout our work, we actually compensate for the increase in average energy per transmitted symbol $x_{v,n}$ (or equivalently $x_{v,q}$) by decreasing the average energy per symbol $x_n$, i.e., making the constellation more compact. This results in a reduced minimum Euclidean distance between signal points, and thus some error performance degradation at the receiver output. In the case of BPSK, the symbols $x_n \in \{\pm \sqrt{E_b} \}$ are simply replaced by the symbols $x_n \in \{\pm \sqrt{E_b}/\omega \}$, where $\omega = 1 + \delta \cdot (C^2 - 1)$. It is easy to show that extending $K$ symbols by a factor $C$ and leaving $(N - K)$ of them unchanged in a frame of $N$ BPSK symbols $x_n \in \{\pm \sqrt{E_b}/\omega \}$ leads to an average energy per symbol equal to $E_b$, and thus identical to the energy per symbol obtained using classical SLM with $x_n \in \{\pm \sqrt{E_b} \}$. If the transmission channel over each BPSK subcarrier can be modeled as AWGN, the bit error probability $p_{eb}$ obtained with the proposed SLM method in the absence of any SI detection error can then be expressed as

$$p_{eb} = \frac{1 - \delta}{2} \cdot \text{erfc} \left( \frac{1}{\omega} \cdot \frac{E_b}{N_0} \right) + \frac{\delta}{2} \cdot \text{erfc} \left( \frac{C^2 \cdot E_b}{\omega \cdot N_0} \right).$$  

(12)

This equation indicates that, at high SNR, the performance gap over AWGN channel between our SLM method and SLM-PSI is equal to $10 \cdot \log_{10}[\omega]$ dB, which corresponds exactly to the energy increase $G$ given by (2). However, as shown by (4), the transmission channel over each subcarrier can actually be modeled more accurately as flat Rayleigh fading rather than AWGN. Based on (12), we can demonstrate that the bit error probability for BPSK over flat Rayleigh fading channel can be approximated at high SNR using [7]:

$$p_{eb} \approx \frac{1 + \delta \cdot (C^2 - 1) \cdot (C^2 - \delta \cdot (C^2 - 1))}{4C^2} \cdot \left( \frac{E_b}{N_0} \right)^{-1}.$$  

(13)

Using (13), it is possible to evaluate the expected error performance gap at high SNR between our SLM technique and SLM-PSI in the absence of any erroneous SI detection event. Note that the bit error probability expression for SLM-PSI is obtained by computing (13) with $C = 1$ and $\delta = 0$. When $C = 1.2$ and $\delta = \frac{1}{2}$, we find a gap of approximately 0.14 dB, which is indeed significantly smaller than the energy increase $G = 0.70$ dB. This theoretical result is therefore coherent with the observations made from the BER plots in Fig. 3.

D. PAPR reduction performance

Due to the lack of space, we are not able to display any result on the PAPR reduction achieved with the proposed SLM technique. However, various computer simulations have indicated that the complementary cumulative distribution function (CCDF) of the PAPR obtained with our SLM method is strictly identical to that of classical SLM [7].

IV. CONCLUSION

We have proposed a simple SLM technique for PAPR reduction in OFDM that does not require the explicit transmission of SI bits. Our investigations have shown that this technique is particularly attractive for systems using a large number of subcarriers. In fact, the probability of SI detection error can be made very small by increasing the extension factor and/or the number of subcarriers. In cases where this probability becomes sufficiently low, the BER performance difference between the proposed technique and classical SLM using error-free side information has been shown to be negligible.

REFERENCES


